



RADISH: Sound and Complete Race Detection in Software and Hardware

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Uses of Race Detection

multithreaded
record+replay

simplifying
consistency
models

atomicity
checking

atomicity
enforcement

concurrency
bug detection

testing &
verification

determinism
checking

determinism
enforcement

many uses require
sound+complete detection

software

Flanagan and Freund, PLDI 2009

hardware

Min and Choi, ASPLOS 1991

Muzahid et al., ISCA 2009

Prvulovic, HPCA 2006

complementary strengths



sound+complete
static analysis



\$ coherence
event-based



slow
polling-based

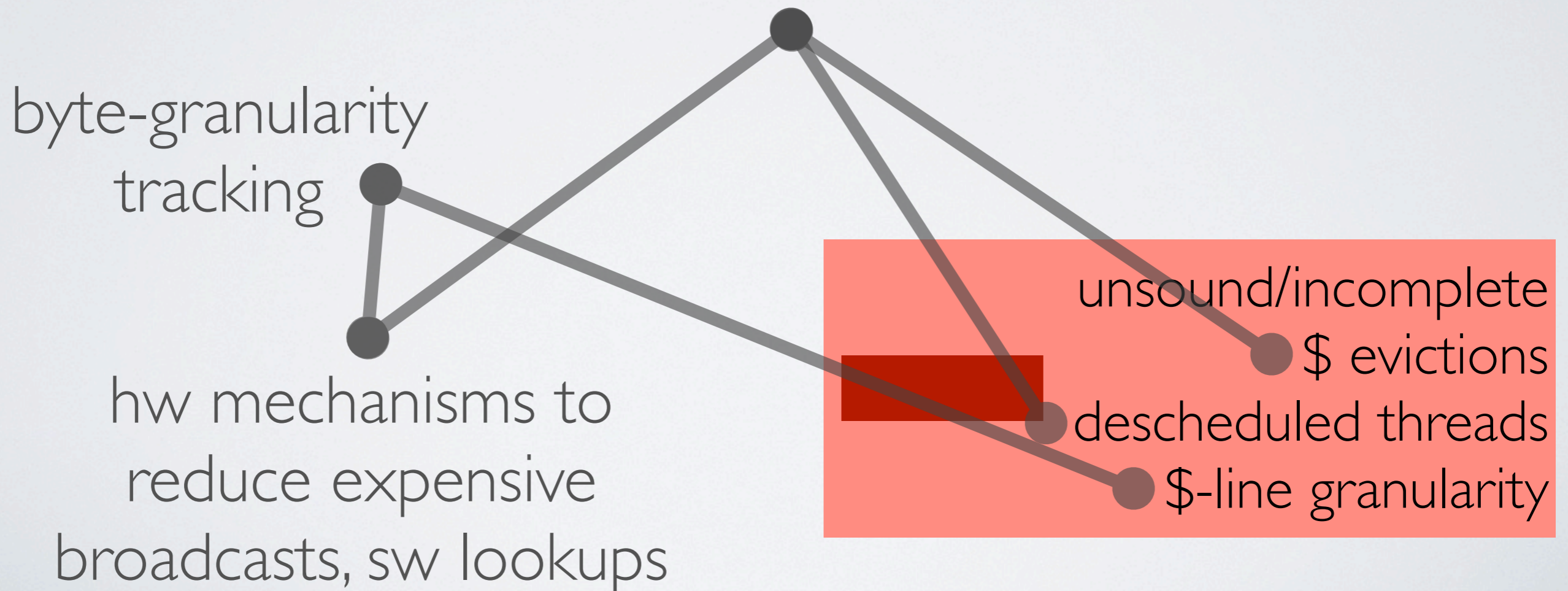


unsound/incomplete
\$ evictions
descheduled threads
\$-line granularity

RADISH overview

sound+complete race detection in sw+hw

use sw to virtualize hw
resources via “revision control”



outline

in-\$ RADISH

results

happens-before
data race detection

full RADISH

conclusions


data races

Lamport, CACM 1978

2 **concurrent** accesses to the same memory location,
 ≥ 1 of which is a write



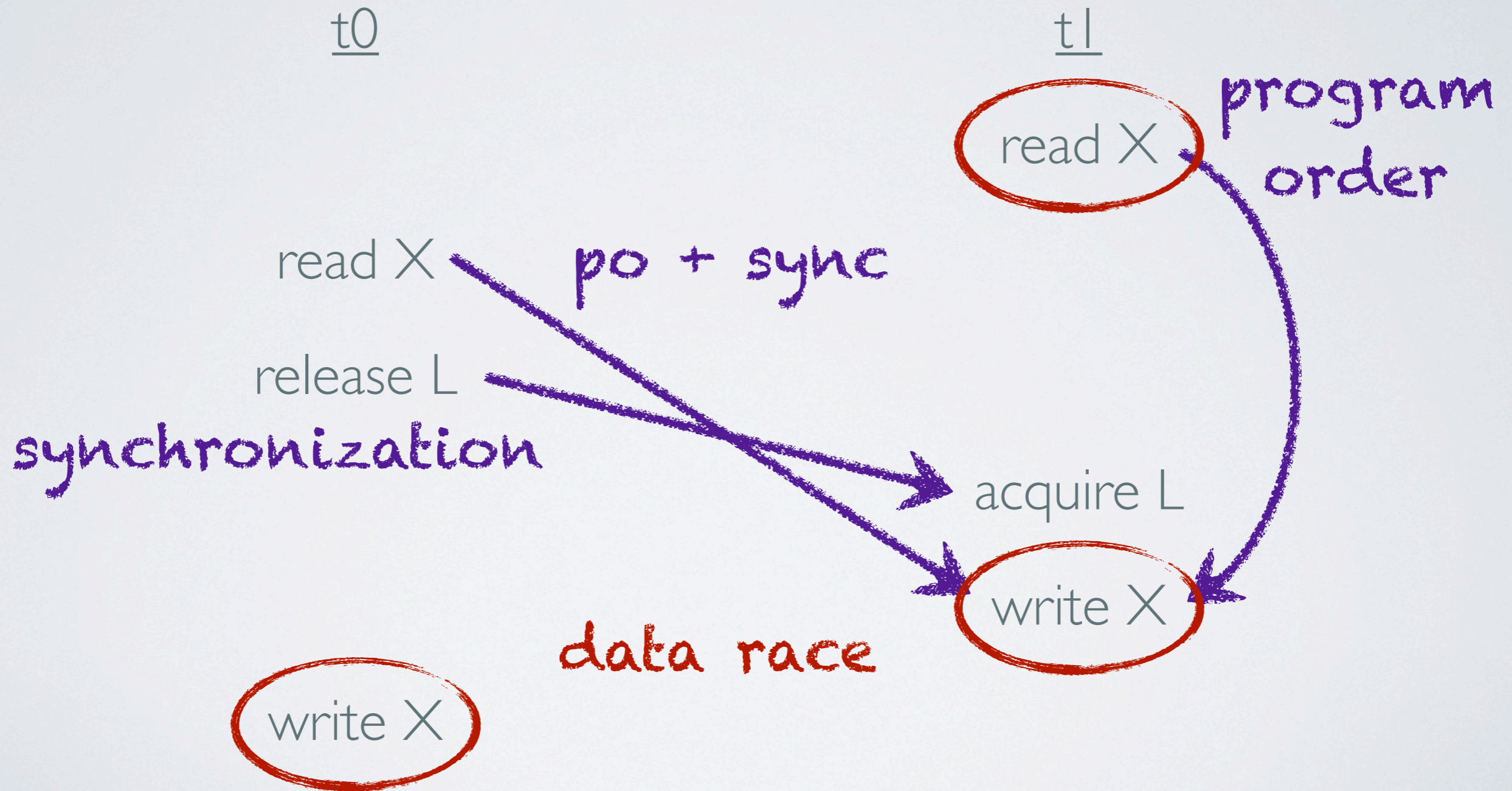
unordered wrt the **happens-before relation**



transitive closure of
program order + synchronization order

data races

Lamport, CACM 1978



happens-before race detection

canonical sound+complete approach

Fidge, Computer 1991 Mattern, IWPDPA 1989

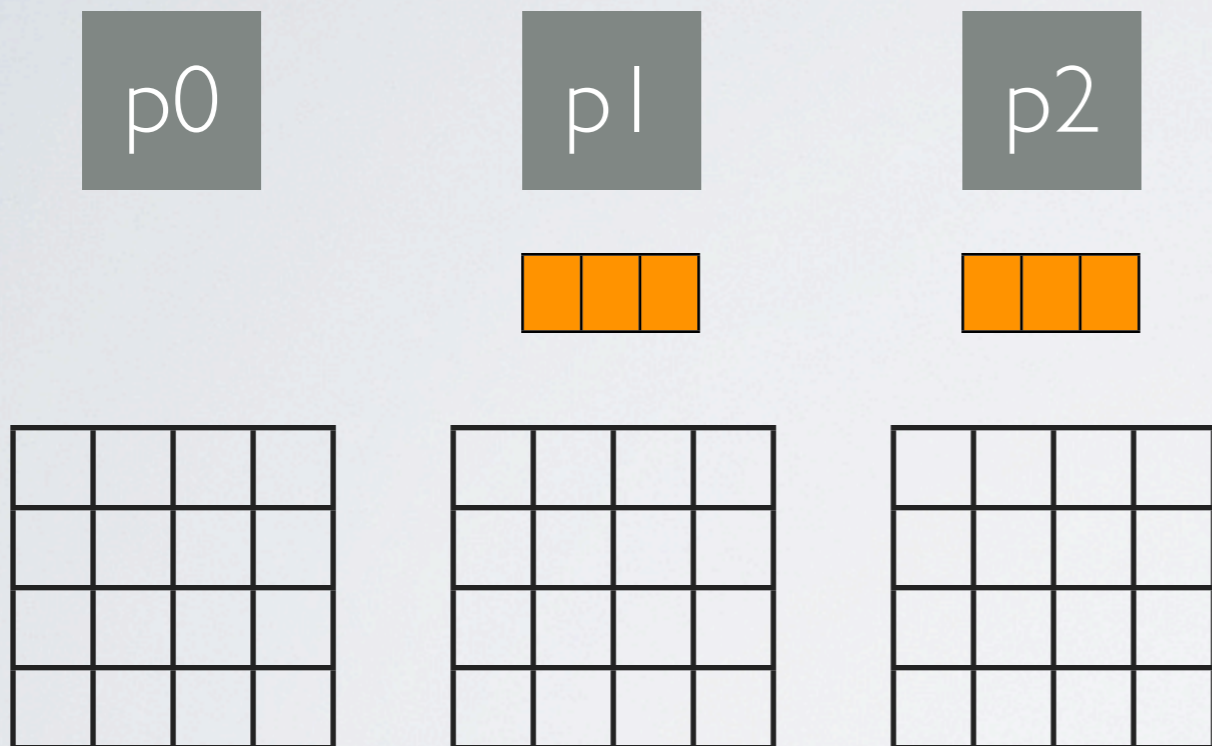
per-**thread** metadata
read ordered with last write

thread	last synchronized with
t0	t1@T, t2@U

write ordered with last write
and all last reads
per-**location** metadata

location	last write	last reads
X	t2@T	t0@U, t1@-, t2@W

mapping to hardware



unbounded # threads

unbounded # locations

thread	last synchronized with
t0	t1 @ [cache] @ U

location	last write	last reads
X	t2 @ T	t0 @ U, t1 @ -, t2 @ W

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strawman

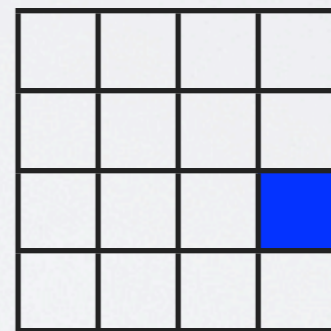
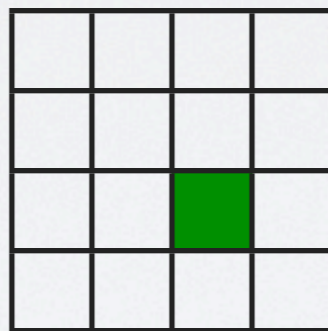
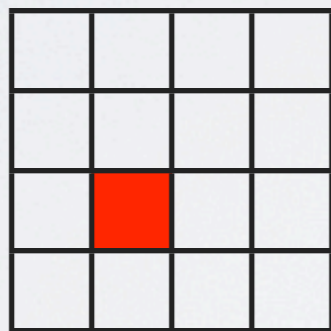
all metadata is in hv, so broadcast on every access?

p0

p1

p2

p0



write X

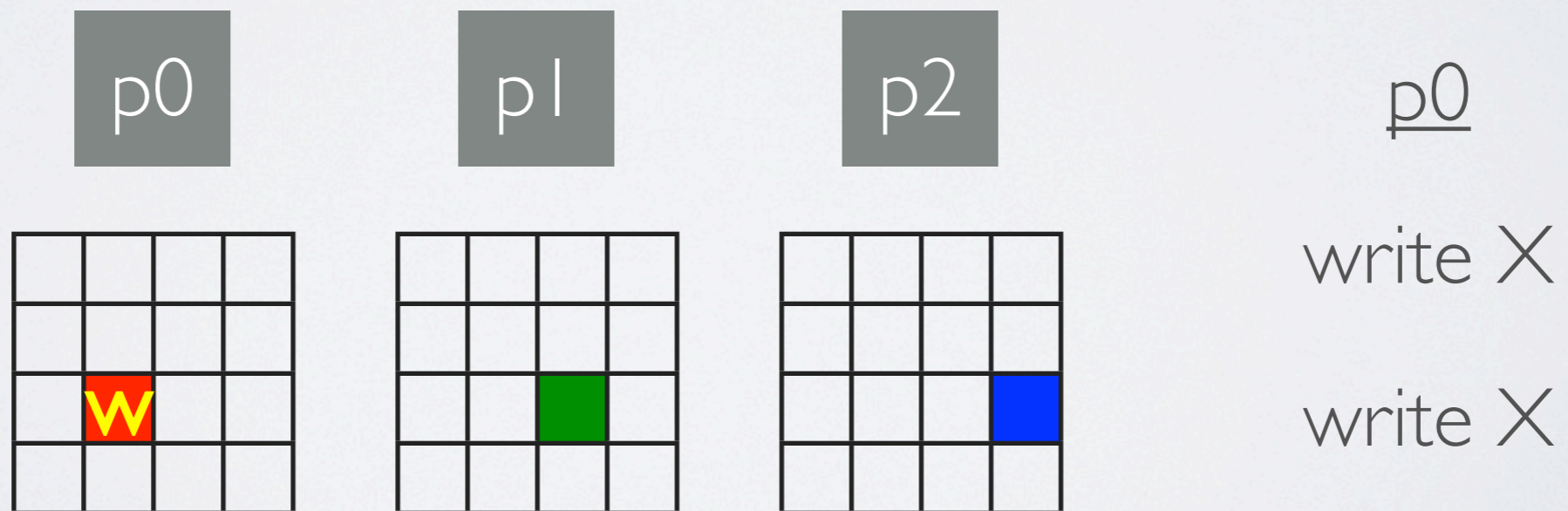
write X

local permissions cache what can be done without communication

local permissions

READ, WRITE or NONE permissions
to each byte in a \$ line

updated only on permissions violations
and coherence events



outline

in-\$ RADISH

results

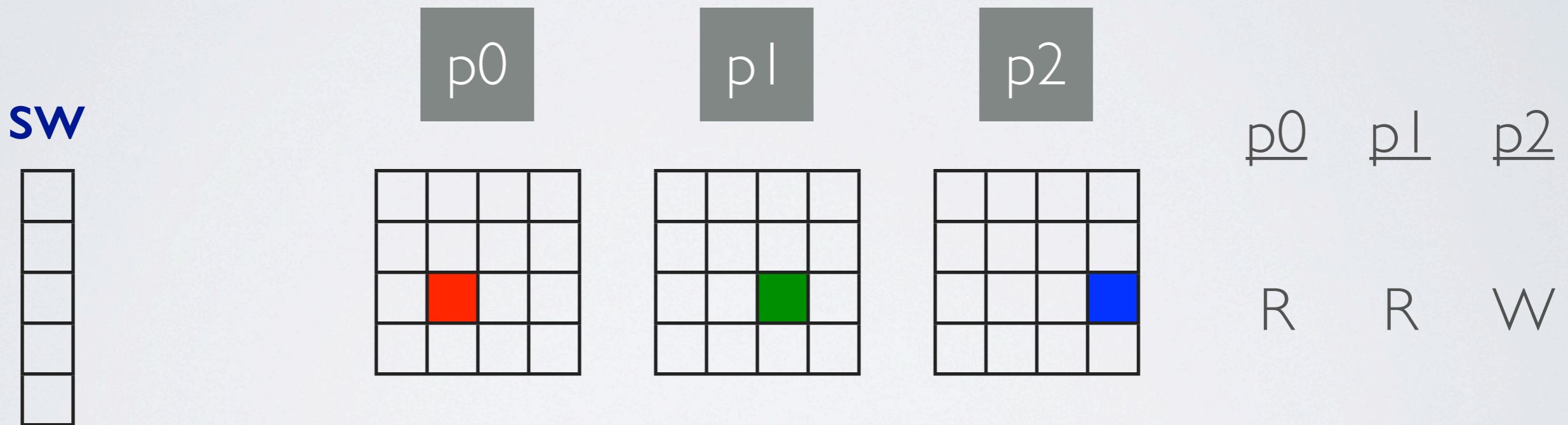
happens-before
data race detection

full RADISH

conclusions

strawman

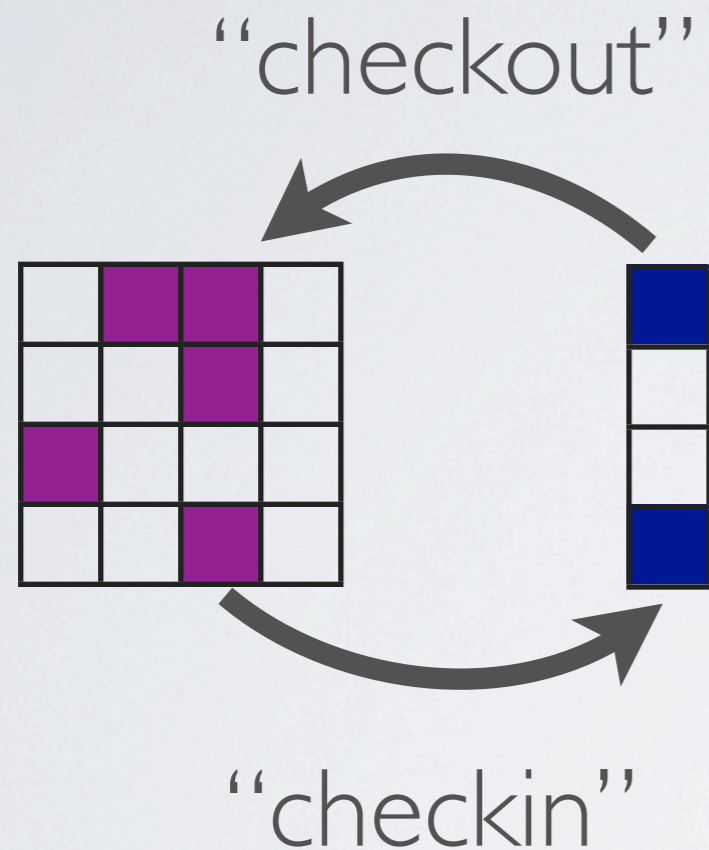
metadata can be in hw or sw
so check sw on every access?



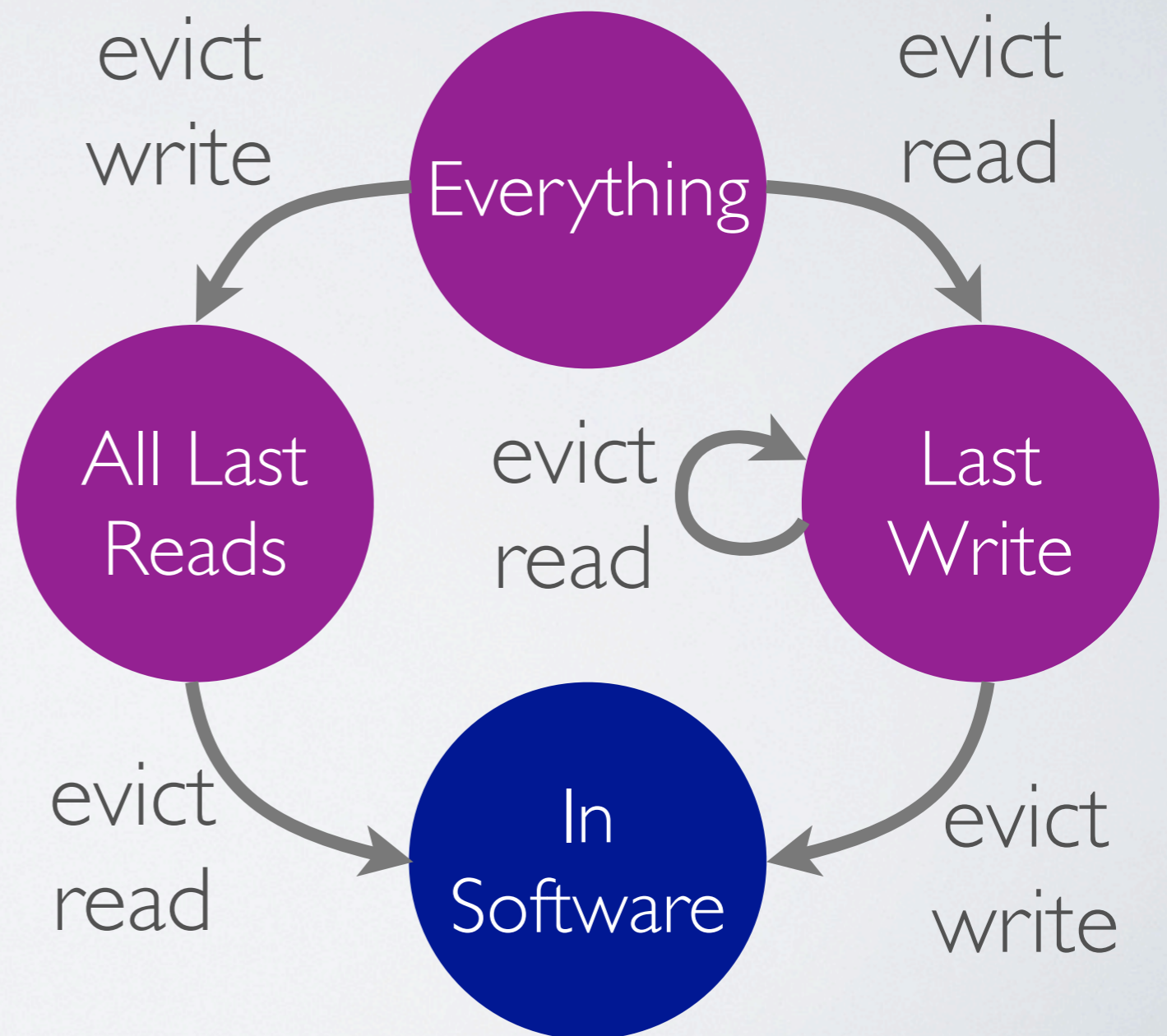
in-hardware status summarizes
what metadata resides in hw

in-hardware status

what can we figure out without going to sw?



set IHS on checkout,
degrade on \$ evictions



also in the paper

leveraging type-safe languages to
reduce metadata space overheads

asynchronous software lookups to
reduce overheads

outline

in-\$ RADISH

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data race detection

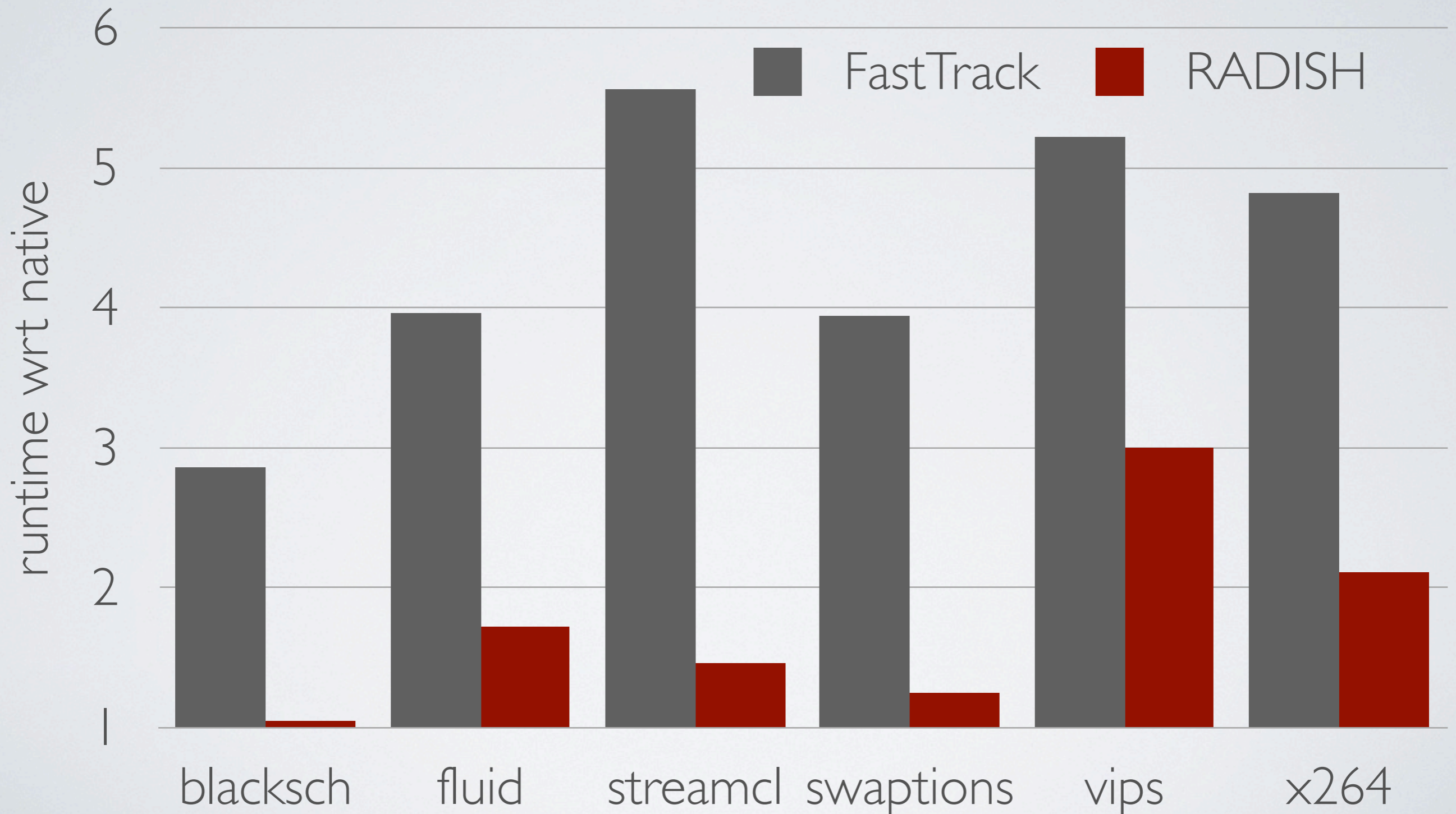
full RADISH

conclusions

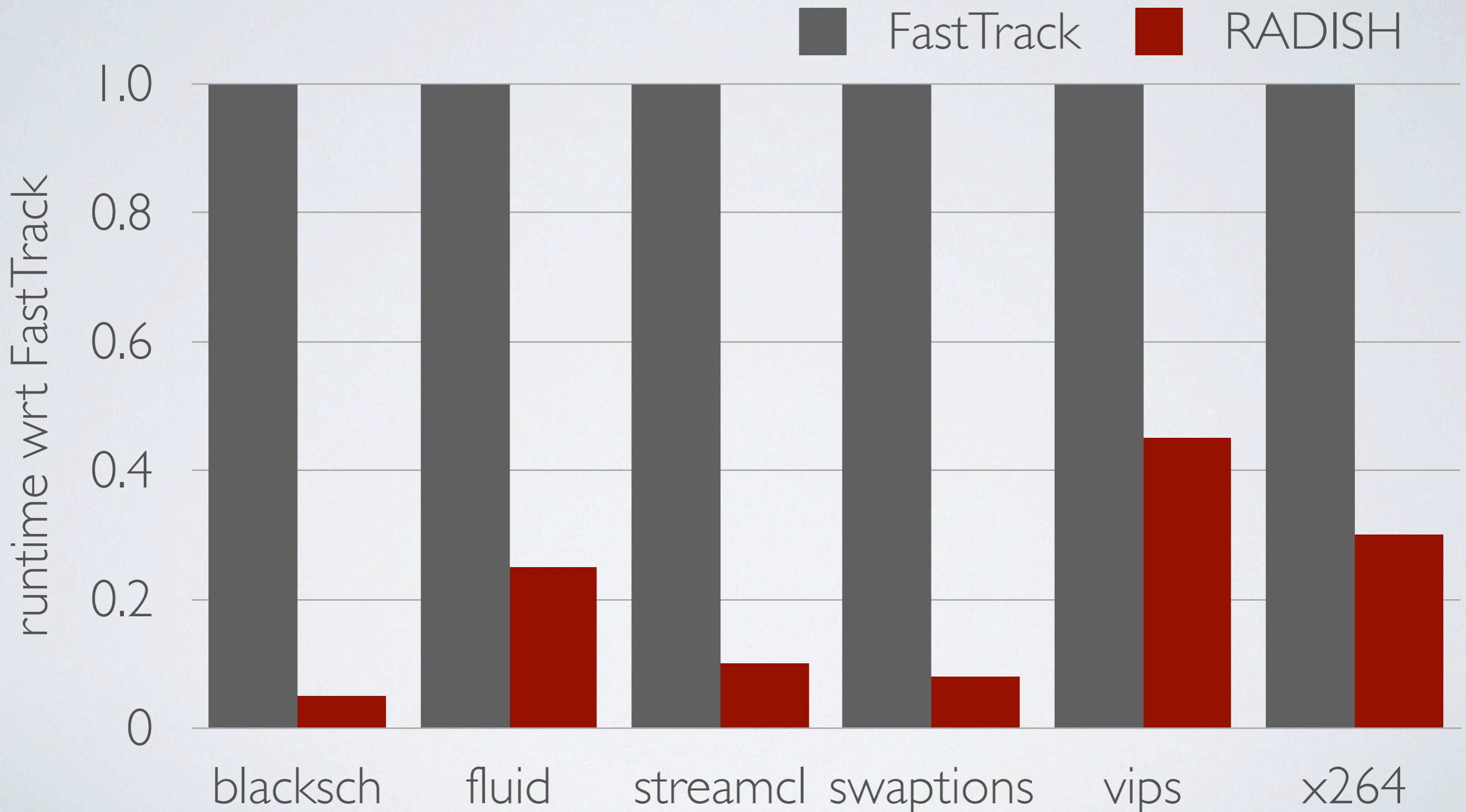
simulation methodology

- Pin-based simulator
- 8 cores, MESI coherence
- 8-way 64KB L1, 8-way 256KB private L2, 16-way 16MB L3
- PARSEC 2.1
- compare with FastTrack [Flanagan and Freund, PLDI 2009]

runtime compared to native



runtime compared to FastTrack



conclusions

sound+complete race detection in hw+sw

unmodified cache design

much faster than software-only race detection



thanks!