CS-XXX: Graduate Programming Languages

Lecture 25 — Multiple Inheritance and Interfaces

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Multiple Inheritance

Why not allow class C extends C1,C2,...{...} (and C ≤ C1 and C ≤ C2)?

What everyone agrees: C++ has it and Java doesn’t

All we’ll do: Understand some basic problems it introduces and how interfaces get most of the benefits and some of the problems

Problem sources:
- Class hierarchy is a dag, not a tree (not true with interfaces)
- Subtype hierarchy is a dag, not a tree (true with interfaces)
Diamond Issues

If \( C \) extends \( C_1 \) and \( C_2 \) and \( C_1, C_2 \) have a common superclass \( D \) (perhaps transitively), our class hierarchy has a diamond

- If \( D \) has a field \( f \), should \( C \) have one field \( f \) or two?
- If \( D \) has a method \( m \), \( C_1 \) and \( C_2 \) will have a clash
- If subsumption is coercive (changing method-lookup), how we subsume from \( C \) to \( D \) affects run-time behavior (incoherent)

Diamonds are common, largely because of types like \texttt{Object} with methods like \texttt{equals}
Multiple Inheritance, Method-Name Clash

If $C$ extends $C_1$ and $C_2$, which both define a method $m$, what does $C$ mean?

Possibilities:

1. Reject declaration of $C$ (Too restrictive with diamonds)
2. Require $C$ to override $m$ (Possibly with directed resends)
3. “Left-side” ($C_1$) wins (Must decide if upcast to “right-side” ($C_2$) coerces to use $C_2$’s $m$ or not)
4. $C$ gets both methods (Now upcasts definitely coercive and with diamonds we lose coherence)
5. Other?
Implementation Issues

This isn’t an implementation course, but many semantic issues regarding multiple inheritance have been heavily influenced by clever implementations

▶ In particular, accessing members of self via compile-time offsets...

▶ ... which won’t work with multiple inheritance unless upcasts “adjust” the self pointer

That’s one reason C++ has different kinds of casts

Better to think semantically first (how should subsumption affect the behavior of method-lookup) and implementation-wise second (what can I optimize based on the class/type hierarchy)
Digression: Casts

A “cast” can mean many things (cf. C++).

At the language level:
- **upcast**: no run-time effect until we get to static overloading
- **downcast**: run-time failure or no-effect
- **conversion**: key question is round-tripping
- “reinterpret bits”: not well-defined

At the implementation level:
- **upcast**: usually no run-time effect but see last slide
- **downcast**: usually only run-time effect is failure, but...
- **conversion**: same as at language level
- “reinterpret bits”: no effect (by definition)
Least Supertypes

Consider if $e_1$ then $e_2$ else $e_3$ (or in C++/Java, $e_1 \ ? \ e_2 : e_3$)

- We know $e_2$ and $e_3$ must have the same type

With subtyping, they just need a common supertype

- Should pick the least (most-specific) type
- Single inheritance: the closest common ancestor in the class-hierarchy tree
- Multiple inheritance: there may be no least common supertype

Example: $C_1$ extends $D_1, D_2$ and $C_2$ extends $D_1, D_2$

Solutions: Reject (i.e., programmer must insert explicit casts to pick a common supertype)
Multiple Inheritance Summary

- Method clashes (what does inheriting \( m \) mean)
- Diamond issues (coherence issues, shared (?) fields)
- Implementation issues (slower method-lookup)
- Least supertypes (may be ambiguous)

Complicated constructs lead to difficult language design
- Doesn’t necessarily mean they are bad ideas

Now discuss *interfaces* and see how (and how not) multiple interfaces are simpler than multiple inheritance...
Interfaces

An interface is just a (named) (object) type. Example:

```java
interface I { Int get_x(); Bool compare(I); }
```

A class can implement an interface. Example:

```java
class C implements I {
    Int x;
    Int get_x() {x}
    Bool compare(I i) {...} // note argument type
}
```

If $C$ implements $I$, then $C \leq I$

Requiring explicit “implements” hinders extensibility, but simplifies type-checking (a little)

Basically, $C$ implements $I$ if $C$ could extend a class with all abstract methods from $I$
Interfaces, continued

Subinterfaces (interface J extends I { ...}) work exactly as subtyping suggests they should

An unnecessary addition to a language with abstract classes and multiple inheritance, but what about single inheritance and multiple interfaces:

\[
\text{class C extends D implements I1,I2,...,In}
\]

- Method clashes (no problem, inherit from D)
- Diamond issues (no problem, no implementation diamond)
- Implementation issues (still a “problem”, different object of type I will have different layouts)
- Least supertypes (still a problem, this is a typing issue)
Using Interfaces

Although it requires more keystrokes and makes efficient implementation harder, it may make sense (be more extensible) to:

- Use interface types for all fields and variables
- Don’t use constructors directly: For class $C$ implementing $I$, write:
  
  ```
  I makeI(...) { new C(...) }
  ```

This is related to “factory patterns”; constructors are behind a level of indirection

It is using named object-types instead of class-based types