

Problem Set 1

599i: Communication Complexity

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Note: This is an optional homework assignment; however, I will grade it for those who choose to submit. Please send your submissions by email.

Problem 1. Suppose Alice gets a string $x \in \{0,1\}^n$ that has more 0s than 1s, and Bob gets a string $y \in \{0,1\}^n$ that has more 1s than 0s. They wish to communicate to find a coordinate i where $x_i \neq y_i$. Show that at least $2 \log n$ bits of communication are required.

Problem 2. Let $f : \{0,1\}^n \times \{0,1\}^n \rightarrow \{0,1\}$ be a function with the property that for every $x \in \{0,1\}^n$, there are exactly t choices for y with $f(x,y) = 1$, and for every $y \in \{0,1\}^n$, there are exactly t choices for x with $f(x,y) = 1$. Our goal is to show that f must admit a 0-cover of size at most $(2n)^t$.

To do this, we use Hall's matching theorem. Let L, R be disjoint sets of the same size. Given a bipartite graph with vertex set $L \cup R$, where every edge belongs to the set $L \times R$, Hall's theorem says that if every subset $S \subseteq L$ is connected to at least $|S|$ vertices in R , then the graph must contain a matching, namely $|L|$ disjoint edges.

- Use Hall's theorem to prove that there are permutations π_1, \dots, π_t such that $f(x,y) = 1$ if and only if $y = \pi_r(x)$ for some r .
- Give a 0-cover of size $(2n)^t$ for f .

Problem 3. The matrices we are working with have 0/1 entries, so one can view these entries as coming from any field – for instance, we can view them as real numbers, rational numbers, or elements of the finite field \mathbb{F}_2 . This is important because the value of the rank may depend on the field used.

1. Give an example of a 3×3 matrix whose rank over \mathbb{F}_2 is not the same as its rank over the reals.
2. Prove that if M has 0/1 entries, then the rank of M over the reals is equal to its rank over the rationals, which is at least as large as its rank over \mathbb{F}_2 .
3. For every r , show that there is a matrix whose \mathbb{F}_2 -rank and communication complexity are both r . Conclude that the log-rank conjecture is false over \mathbb{F}_2 .

Problem 4. Show that there is a Boolean matrix of rank at most $2r$ with 2^r distinct rows and 2^r distinct columns.

Problem 5. Prove that John's theorem is tight.