Threads & Synchronization

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Outline

- Previous lectures: concurrency
  - threads vs. processes
  - how to implement threads?
- Next few lectures:
  - how to write multithreaded programs?
    - Main challenge: how to eliminate race conditions? how to synchronize?
    - Solutions: locks, semaphore, conditional variables, monitors, ...

Independent vs. Cooperating

- Independent threads
  - **no state shared with other threads**
  - deterministic --- input state determines result
  - reproducible
  - scheduling order does not matter

- Cooperating threads
  - **shared state**
  - non-deterministic
  - non-reproducible

Non-reproducibility and non-determinism means that bugs can be intermittent. This makes debugging really hard.

Why allow cooperating threads ?

Computer programs at some level have to cooperate

- Share resources/information
  - one computer many users/programs
  - one bank balance, many ATMs
- Speedup
  - overlap IO and computation
  - multiprocessors -- chop up programs into smaller pieces
- Modularity
  - chop large problem up into simpler pieces
  - For example: "delatex foo.tex | spell | sort | uniq | wc"

Example: Shared counter

- Yahoo gets millions of hits a day. Uses multiple threads (on multiple processors) to speed things up.
- Simple shared state error: each thread increments a shared counter to track the number of hits today:

```
...
hits = hits + 1;
...
```

- What happens when two threads execute this code concurrently?

Problem with shared counters

- One possible result: lost update!

hits = 0

time

T1

read hits (0)

hits = 0 + 1

T2

read hits (0)

hits = 0 + 1

hits = 1

- **One other possible result: everything works.**
  - Bugs are frequently intermittent. Makes debugging hard.
  - This is called “race condition”

Race conditions

- Race condition: timing dependent error involving shared state.
  - whether it happens depends on how threads are scheduled
- *Hard* because:
  - ***must make sure all possible schedules are safe***. Number of possible schedule permutations is huge.

```
if (n == stack_size)   /* A */
        return full;   /* B */
stack[n] = v;          /* C */
n = n + 1;             /* D */
```

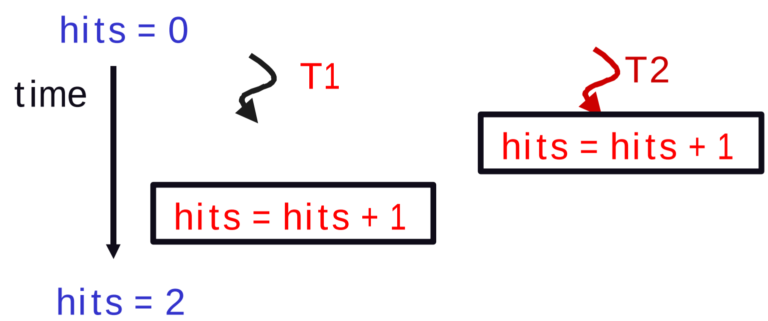
Stack Race Conditions

```
if (n == stack_size)   /* A */      if (n == stack_size)   /* A' */
        return full;   /* B */              return full;   /* B' */
stack[n] = v;          /* C */      stack[n] = v;          /* C' */
n = n + 1;             /* D */      n = n + 1;             /* D' */
```

- Some bad schedules:
  - AA'CC'DD' → overwrites
  - ACA'DC'D' → overflow
  - How many???
- Bugs are intermittent. Timing dependent = small changes (adding a print stmt, different machine) can hide bug.

Preventing race conditions: atomicity

- atomic unit = instruction sequence guaranteed to execute indivisibly (also, a “critical section”).
    - If two threads execute the same atomic unit at the same time, one thread will execute the whole sequence before the other begins.



- How to make multiple instructions seem like one atomic one?

A few definitions

- Critical section:
    - piece of code that only one thread can execute at once. Only one thread at a time will get into the section of code.
- Mutual exclusion:
    - ensuring that only one thread does a particular thing at a time. One thread doing it excludes the other, and vice versa.
- Lock: prevents someone from doing something
    - lock before entering critical section, before accessing shared data
    - unlock when leaving, after done accessing shared data
    - wait if locked
- Synchronization:
    - using atomic operations to ensure cooperation between threads

Example: the Too-Much-Milk problem

- Consider a bunch of roommates in a house

```
Person() {

    while (1)  {
        Dosomething();

        if (!CheckMilk)
            BuyMilk();
    }

}
```

Goal: 1. never more than one person buys
2. someone buys if needed (otherwise "starvation")

Example: the Too-Much-Milk problem

	Person A	**Person B**
3:00	Look in fridge. Out of milk	
3:05	Leave for store	
3:10	Arrive at store	Look in fridge. Out of milk
3:15	Buy milk	Leave for store
3:20	Arrive home, put milk away	Arrive at store
3:25		Buy milk
3:30		Arrive home, put milk away
		Oh no !

Too much milk: solution #1

- Basic idea:
  - leave a note (kind of like "lock")
  - remove note (kind of like "unlock")
  - don't buy if there is a note (wait)

```
if (noMilk) {
  if (noNote) {
      leave Note;
      buy milk;
      remove Note
  }
}
```

Why solution #1 does not work ?

	Thread A	Thread B
3:00	if (noMilk) {	
3:05	if (noNote) {	
3:10		if (noMilk) {
3:15		if (noNote) {
3:20	leave Note;	leave Note;
3:25	buy milk;	buy milk;
3:30	remove Note } }	remove Note } }

Threads can get context-switched at any time !

Too much milk: solution #2

Thread A

```
leave NoteA
if (noNoteB) {
    if (noMilk)
        buy milk
}
remove NoteA
```

Thread B

```
leave NoteB
if (noNoteA) {
    if (noMilk)
        buy milk
}
remove NoteB
```

Problem with Solution #2

Thread A	**Thread B**
leave NoteA	
	leave NoteB
if (noNoteB) {	
if (noMilk)	
buy milk	
}	
	if (noNoteA) {
	if (noMilk)
	buy milk
	}
remove NoteA	
	remove NoteB

Problem: neither thread to buy milk --- think other is going to buy --- ***starvation!***

Too much milk: solution #3

Thread A	Thread B

Thread A

```
leave NoteA
while (NoteB)
      do nothing;
if (noMilk)
      buy milk;
remove NoteA
```

Thread B

```
leave NoteB
if (noNoteA) {
      if (noMilk)
            buy milk;
}
remove NoteB
```

Either safe for me to buy or others will buy !

Solution #3: a scenario

Thread A

```
leave NoteA

while (NoteB)
     do nothing;






if (noMilk)
     buy milk;
remove NoteA
```

Thread B

```


leave NoteB



if (noNoteA) {
     if (noMilk)
           buy milk;
}
remove NoteB
```

Solution #3: another scenario

Thread A

```
leave NoteA
while (NoteB)
    do nothing;

if (noMilk)
    buy milk;
remove NoteA
```

Thread B

```
leave NoteB
if (noNoteA) {
    if (noMilk)
        buy milk;
}
remove NoteB
```

Solution #3: another scenario

Thread A

```
leave NoteA
while (NoteB)
    do nothing;
if (noMilk)

    buy milk;
remove NoteA
```

Thread B

```
leave NoteB
if (noNoteA) {
    if (noMilk)
        buy milk;
}
remove NoteB
```

Question: any criticisms on this style of providing mutual exclusion?

Locks using load/store

- Dekker's algorithm, later simplified by Peterson
- No hardware support required

```
lockedA = true;
turn = B;
while (lockedB && turn != A);
<critical section>
lockedA = false;
```

```
lockedB = true;
turn = A;
while (lockedA && turn != B);
<critical section>
lockedB = false;
```

Scenario 1

```
lockedA = true;
turn = B;
while (lockedB && turn != A);
<critical section>



lockedA = false;
```

```



lockedB = true;
turn = A;
while (lockedA && turn != B);
        <blocks>

        <released>
<critical section>
```

Scenario 2

```
lockedA = true;
turn = B;
while (lockedB && turn != A);
<critical section>
lockedA = false;
```

```
lockedB = true;
turn = A;
while (lockedA && turn != B);
    <blocks>
```

Scenario 3

```
lockedA = true;
turn = B;
while (lockedB && turn != A);
    <blocks>
```

```
lockedB = true;
turn = A;
while (lockedA && turn != B);
<critical section>
lockedB = false;
```

A better solution

- Have hardware provide better primitives than simple load and store.
- Build higher-level programming abstractions on this new hardware support.
- Example: using locks as an atomic building block

Lock::Acquire --- wait until lock is free, then grabs it

Lock::Release --- unlock, waking up a waiter if any

These must be atomic operations --- if two threads are waiting for the lock, and both see it is free, only one grabs it!

```
lock -> Acquire();
if (nomilk)
    buy milk;
lock -> Release();
```

Announcements

- Assignment 1 will be online tonight:
  - Send us groupings by Wednesday/thursday
  - Design due by next Tuesday