

Due: May 12, 2021

Problem 1: Strong subadditivity of the quantum entropy

We have established that the quantum relative entropy is monotone decreasing under partial trace: For bipartite states $\rho, \sigma \in \mathcal{D}(\mathcal{H}_1 \otimes \mathcal{H}_2)$, it holds that

$$\mathcal{S}(\text{Tr}_{\mathcal{H}_2}(\rho) \parallel \text{Tr}_{\mathcal{H}_2}(\sigma)) \leq \mathcal{S}(\rho \parallel \sigma). \quad (0.1)$$

Let $\rho^{ABC} \in \mathcal{D}(\mathcal{H}_A \otimes \mathcal{H}_B \otimes \mathcal{H}_C)$ be a tripartite state, and denote the partial traces

$$\rho^{AB} = \text{Tr}_C(\rho^{ABC}), \quad \rho^{BC} = \text{Tr}_A(\rho^{ABC}), \quad \rho^A = \text{Tr}_{BC}(\rho^{ABC}), \quad \rho^B = \text{Tr}_{AC}(\rho^{ABC}).$$

Using (0.1) with the states $\rho = \rho^{ABC}$ and $\sigma = \rho^A \otimes \rho^{BC}$, prove that strong subadditivity holds:

$$\mathcal{S}(\rho^{ABC}) + \mathcal{S}(\rho^B) \leq \mathcal{S}(\rho^{AB}) + \mathcal{S}(\rho^{BC}), \quad (0.2)$$

where \mathcal{S} denotes the von Neumann entropy.

Remark 0.1. Recall from Lecture 9 that (0.2) can be interpreted as establishing nonnegativity of the quantum conditional mutual information:

$$\mathbf{I}(A, C \mid B)_\rho := \mathcal{S}(\rho^{AB}) + \mathcal{S}(\rho^{BC}) - \mathcal{S}(\rho^{ABC}) - \mathcal{S}(\rho^B).$$

Problem 2: Monotonicity of relative entropy under quantum channels

Suppose that $\mathcal{E} : \mathbb{M}_n(\mathbb{C}) \rightarrow \mathbb{M}_k(\mathbb{C})$ is a quantum channel (i.e., a CPT map). Use Theorem 1.3 in Lecture 10, along with (0.1), to prove that for all $\rho, \sigma \in \mathcal{D}(\mathbb{M}_n(\mathbb{C}))$,

$$\mathcal{S}(\mathcal{E}(\rho) \parallel \mathcal{E}(\sigma)) \leq \mathcal{S}(\rho \parallel \sigma).$$

Problem 3: Quantum Pinsker inequality

Here you will prove the quantum Pinsker inequality, that for any density matrices ρ, σ , we have

$$\mathcal{S}(\rho \parallel \sigma) \geq \frac{1}{2} \|\rho - \sigma\|_1^2, \quad (0.3)$$

where $\|\cdot\|_1$ is the Schatten 1-norm:

$$\|A\|_1 = \text{Tr}(|A|) = \text{Tr}(\sqrt{A^*A}).$$

You can assume the classical Pinsker inequality for probability distributions p, q which asserts that

$$D(p \parallel q) \geq \frac{1}{2} \|p - q\|_1^2,$$

and here $\|\cdot\|_1$ is the standard ℓ_1 norm.

1. Prove that for any density matrices ρ and σ , it holds that $\mathcal{S}(D_\rho \parallel D_\sigma) \leq \mathcal{S}(\rho \parallel \sigma)$, where D_ρ and D_σ are ρ and σ with the off-diagonal entries zeroed out. [Hint: It may help to recall Problem 2(a) on the first homework.]
2. Prove (0.3). [Hint: First reduce to the case when $\rho - \sigma$ is diagonal.]

Problem 4: Squashed entanglement

Denote a bipartite system $\mathcal{H}_A \otimes \mathcal{H}_C$. For a given state $\rho^{AC} \in \mathcal{D}(\mathcal{H}_A \otimes \mathcal{H}_C)$, consider all extensions $\rho^{ABC} \in \mathcal{D}(\mathcal{H}_A \otimes \mathcal{H}_B \otimes \mathcal{H}_C)$ for some auxiliary Hilbert space \mathcal{H}_B such that $\rho^{AC} = \text{Tr}_B(\rho^{ABC})$. Define the *squashed entanglement* of ρ^{AC} by

$$\mathbf{E}^{\text{sq}}(\rho^{AC}) := \frac{1}{2} \inf_{\rho^{ABC}} (\mathcal{S}(\rho^{AB}) + \mathcal{S}(\rho^{BC}) - \mathcal{S}(\rho^{ABC}) - \mathcal{S}(\rho^B)),$$

where the infimum is over extensions ρ^{ABC} of ρ^{AC} (in particular, over all finite-dimensional Hilbert spaces \mathcal{H}_B). Note that this quantity is nonnegative by Problem 1.

For intuition, one can use the quantum conditional mutual information (introduced in Lecture 9) to recast this as

$$\mathbf{E}^{\text{sq}}(\rho^{AC}) := \frac{1}{2} \inf_{\rho} \mathbf{I}(A, C | B)_{\rho},$$

where the infimum is again over all extensions ρ of ρ^{AC} . Thus this captures the amount of “mutual quantum information” across the A - C partition that is not explained by some larger environment.

1. Show that if ρ^{AC} is pure, then $\mathbf{E}^{\text{sq}}(\rho^{AC}) = \mathcal{S}(\rho^A)$.
2. An entanglement measure \mathbf{E} is called *faithful* if $\mathbf{E}(\rho^{AC}) = 0$ if and only if ρ^{AC} is separable (across the A - C partition). Show one direction (the easier one) for squashed entanglement: If ρ^{AC} is separable, then $\mathbf{E}^{\text{sq}}(\rho^{AC}) = 0$.

[Remark: It is also true that if ρ^{AC} is not separable, then $\mathbf{E}^{\text{sq}}(\rho^{AC}) > 0$, but this is more difficult, and it was an open question for a while.]

3. Show that \mathbf{E}^{sq} is additive in the sense that

$$\mathbf{E}^{\text{sq}}(\rho^{AC} \otimes \rho^{A'C'}) = \mathbf{E}^{\text{sq}}(\rho^{AC}) + \mathbf{E}^{\text{sq}}(\rho^{A'C'}),$$

where we think of $\rho^{AC} \otimes \rho^{A'C'}$ as a bipartite state in $\mathcal{D}((\mathcal{H}_A \otimes \mathcal{H}_{A'}) \otimes (\mathcal{H}_C \otimes \mathcal{H}_{C'}))$.

For guidance, the proof can be broken into parts:

- (a) Prove that $\mathbf{E}^{\text{sq}}(\rho^{AC} \otimes \rho^{A'C'}) \leq \mathbf{E}^{\text{sq}}(\rho^{AC}) + \mathbf{E}^{\text{sq}}(\rho^{A'C'})$.
- (b) Prove the following chain rule: For any $\sigma^{XYZU} \in \mathcal{D}(\mathcal{H}_X \otimes \mathcal{H}_Y \otimes \mathcal{H}_Z \otimes \mathcal{H}_U)$,

$$\mathbf{I}(XY, Z | U)_{\sigma} = \mathbf{I}(X, Z | U)_{\sigma} + \mathbf{I}(Y, Z | UX)_{\sigma}.$$

- (c) Use the chain rule, along with nonnegativity of the quantum conditional mutual information to show that for *any* state $\rho^{AA'CC'} \in \mathcal{D}((\mathcal{H}_A \otimes \mathcal{H}_{A'}) \otimes (\mathcal{H}_C \otimes \mathcal{H}_{C'}))$, we have

$$\mathbf{E}^{\text{sq}}(\rho^{AA'CC'}) \geq \mathbf{E}^{\text{sq}}(\rho^{AC}) + \mathbf{E}^{\text{sq}}(\rho^{A'C'}).$$