

# Checking Matrix Multiplication

A Randomized Algorithm

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# Problem

- A, B, and C are matrices of size  $n \times n$ .
- Question:

$$A \stackrel{?}{=} BC$$

# Solution 1

$$A \stackrel{?}{=} BC$$

- Perform standard matrix multiplication for  $B \times C$  and compare the results with  $A$ .
- $O(n^3)$  [ $n^2$  output, each output requires  $n$  op]
- State of the art multiplication  $O(n^{2.373})$   
[Redefines the matrix to remove some redundant operations].
- Can we do better?

# Solution 2

$$A \stackrel{?}{=} BC$$

- Choose a vector  $v = (v_1, v_2, \dots, v_n) \in \{-1, 1\}^n$  uniformly at random.
- Check  $Av = (BC)v$ .
- $(BC)v = B(Cv) \leftarrow O(n^2)$ .  
[two matrix-vector multiplication]
- Pretty fast, but how accurate?

# Solution 2

$$A \stackrel{?}{=} BC$$

*Claim : If  $A \neq BC$ , then  $P[Av = BCv] \leq \frac{1}{2}$*

# Solution 2

$$A \stackrel{?}{=} BC$$

- Proof: Let  $D = A - BC$
- Since  $A \succeq BC$ , then  $D \succeq 0$ .
- Let  $i, j$  be indices such that  $D_{ij} \neq 0$ .
- Then:

# Solution 2

$$A \stackrel{?}{=} BC$$

$$D = A - BC$$

$$\begin{bmatrix} D_{i0} & D_{i1} & \dots & D_{ij} & \dots \\ D_{in} & & & & \end{bmatrix} \times \begin{bmatrix} v_0 \\ v_1 \\ \vdots \\ v_j \\ \vdots \\ v_n \end{bmatrix} = \begin{bmatrix} (Dv)_0 \\ (Dv)_1 \\ \vdots \\ (Dv)_i \\ \vdots \\ (Dv)_n \end{bmatrix}$$

$$P[Av = BCv] = P[Dv = 0] \leq P[(Dv)_i = 0]$$

$$(Dv)_i = \sum_{k=1}^n D_{ik} v_k$$

$$(Dv)_i = D_{ij} v_j + \sum_{k \neq j}^n D_{ik} v_k$$

# Questions?

