

Predicate Dispatching: A Unified Theory of Dispatch

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Ernst, ECOOP '98, page 1

Dispatching

Select one *case* from a *generic function*

- object-oriented dispatch, including multi-methods
- classifiers, predicate classes
- pattern matching

Case selection

- applicability
- overriding

Static checking

- completeness
- uniqueness

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Goals

Unify and generalize dispatching mechanisms
Small, orthogonal basic model
Functions are extensible
Static checking

Solution: predicate dispatching

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Outline

Predicate dispatching
Examples
Semantics

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Applicability

Predicates are boolean formulas over class tests
and program expressions

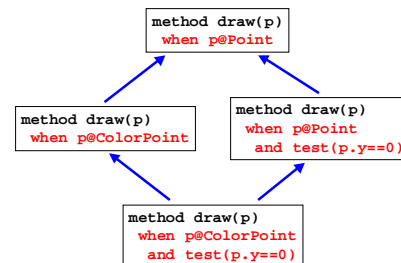
```
pred ::= expr @ class
      | test(expr)
      | let var := expr
      | not pred
      | pred or pred
      | pred and pred
```

```
method draw(p)
  when p@Point and test(p.y == 0)
  { draw point on x axis }
```

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Overriding

Overriding is determined by logical implication



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Object-oriented dispatch (SmallTalk, C++, Java)

Applicability: run-time class is subclass of
specializer

Overriding: subclassing (most specific
specializer)

```
class Point;
method draw(p) when p@Point { ... };

class ColorPoint extends Point;
method draw(p) when p@ColorPoint { ... };

Equivalently: method draw(p@Point) { ... };
              method draw(p@ColorPoint) { ... };
```

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Multi-method dispatch (CLOS, Cecil, Dylan)

Applicability: run-time classes are subclasses
of specializers

Overriding: subclassing over all specializers

```
method equal(p1, p2)
  when p1@Point and p2@Point
  { return p1.x==p2.x and p1.y==p2.y; }

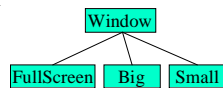
method equal(p1, p2)
  when p1@ColorPoint and p2@ColorPoint
  { return p1.x==p2.x and p1.y==p2.y
    and p1.color==p2.color; }
```

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Classifiers (Kea, Cecil)

Applicability: boolean conditions of runtime state

Overriding: lexical order



```
classify(w@Window)
  as FullScreen when test(w.area() == RootWindow.area())
  as Big when test(w.area() > RootWindow.area()/2)
  as Small otherwise;

method move(w) when w@FullScreen { do nothing }
method move(w) when w@Big { move wireframe outline }
method move(w) when w@Small { move opaque window }
```

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Pattern matching (ML, Haskell)

Applicability: structural equivalence

Overriding: lexical order

```
fun sumList (nil) = 0
  | sumList (h::t) = h + sumList(t);

method sumList(l) when l@Nil
  { return 0; }
method sumList(l)
  when l@Cons and let h:=l.head and let t:=l.tail
  { return h + sumList(t); }
```

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New capability: disjunction

zip: given a pair of lists, return a list of pairs

```
method zip(l1, l2)
  when l1@Nil or l2@Nil
  { return Nil; }

method zip(l1, l2)
  when l1@Cons and l2@Cons
  { return Cons(Pair(l1.head, l2.head),
    zip(l1.tail, l2.tail)); }
```

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Predicate abstractions

Capabilities of in-line predicates:

- return boolean value
- bind variables

```
predicate On_x_axis(p)
  when (p@CartesianPoint and test(p.y == 0))
  or (p@PolarPoint
    and (test(p.theta == 0) or test(p.theta == pi)))

method draw(p) when p@Point { ... }
method draw(p) when On_x_axis(p) { ... }
```

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Run-time behavior

Applicability: evaluate predicates (at run time)
Optimization: common subexpression elimination
Overriding: logical implication
Computed at compile time, used at run time

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Typechecking

Completeness:
Informally, disjunction of predicates = true
Uniqueness: for each pair of predicates,
– disjoint,
– one subsumes the other, or
– their intersection is overridden

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Comparing predicates

Compile-time tautology tests
In general: **undecidable**
Black-box program expressions: **NP-hard**
Small predicate expressions: **fast**
Consistent classes: **complete**

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Future work

More efficient implementation strategies
Separate typechecking

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Contributions of predicate dispatching

Generalizes and subsumes previous techniques in a
common framework
Enables new varieties of tests
Implementation:
<http://www.cs.washington.edu/research/projects/cecil/www/Gud>

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