Planning for Change in a Formal Verification of the Raft Consensus Protocol









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Contributions

First formal proof of Raft's safety *first verified implementation!*

Large-scale Verdi case study stress test; reverification inevitable

Proof engineering lessons affinity lemmas, etc.







Distributed Systems





Reliably deliver procrastination





Also serious infrastructure



One day last summer...

The New York Times

The Stock Market Bell Rings, Computers Fail, Wall Street Cringes

By NATHANIEL POPPER JULY 8, 2015

Problems with technology have at times roiled global financial markets, but the 223-year-old <u>New York Stock</u> <u>Exchange</u> has held itself up as an oasis of humans ready to step in when the computers go haywire.

On Wednesday, however, those working on the trading floor were left helpless when the computer systems at the exchange went down for nearly four hours in the middle of the day, bringing an icon of capitalism's ceaseless energy to a costly halt.

The exchange ultimately returned to action shortly before the closing bell,



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Cyber Sleuths Trac China's Military

The story of a Chinese military staffer's hacking provides a detailed look into Be controlled cyberespionage machinery.



Aw, Snap!





Debt Relief for Students Snarls Market for Their Loans

Federal programs designed to ease the burden of college loans are causing snarls in the bond market and raising concerns that banks may soon ratchet back lending.



The New Bond Market: Algorithms Trump Humans

Computerized trading strategies, or algorithms, are remaking the \$12.7 trillion Treasury market, emulating earlier sea changes in stock and currency trading.



One day last summer...



How distributed systems fail



Related Work



EventML [LADA12, AVoCS15] language for verified distributed systems

IronFleet [SOSP15] liveness, log compaction, serialization

Verdi [PLDI15]

network semantics, transformers, higher-order

Verdi background



Network semantics operational semantics define network behavior

Verified system transformers prove property transfer to adversarial network



 $\forall \Phi S, \operatorname{holds}(\Phi, S, \rightsquigarrow_1) \rightarrow$ holds(transfer(Φ), $T(S), \rightsquigarrow_2$)

Big Picture



Past: Verdi Framework compositional fault tolerance

Present: Verified Raft critical piece of infrastructure

Future:

dynamically upgrading systems program logic



Verification Challenge state machine replication

Raft Algorithm implemented in Verdi

Proof Overview and lessons learned







Replication for fault tolerance





Replication for fault tolerance



Replication correctness



Replication correctness



Internal Correctness

linearizability follows from internal correctness: **state machine safety**

Goal: Verify Raft





Goal: Verify Raft





Verification Challenge state machine replication

Raft Algorithm implemented in Verdi

Proof Overview and lessons learned







Formalizing the network state of the world (P, Σ, T) packets in flight history of I/O data @ nodes

Formalizing the network

 (P, Σ, T)

Formalizing the network

$(P, \Sigma, T) \rightsquigarrow (P', \Sigma', T')$

Defining network semantics

 $\frac{H_{\text{net}}(dst, \ \Sigma[dst], \ src, \ m) = (\sigma', \ o, \ P') \qquad \Sigma' = \Sigma[dst \mapsto \sigma']}{(\{(src, \ dst, \ m)\} \uplus P, \ \Sigma, \ T) \leadsto (P \uplus P', \ \Sigma', \ T + + \langle o \rangle)} \text{ Deliver}$

Defining network semantics

$$\frac{H_{\text{net}}(dst, \ \Sigma[dst], \ src, \ m) = (\sigma', \ o, \ P') \qquad \Sigma' = \Sigma[dst \mapsto \sigma']}{(\{(src, \ dst, \ m)\} \uplus P, \ \Sigma, \ T) \rightsquigarrow (P \uplus P', \ \Sigma', \ T + + \langle o \rangle)} \text{ Deliver}$$

$$\frac{p \in P}{(P, \Sigma, T) \rightsquigarrow (P \uplus \{p\}, \Sigma, T)} \text{ Duplicate}$$

$$\overline{(\{p\} \uplus P, \Sigma, T)} \rightsquigarrow (P, \Sigma, T)$$
 Drop

$$\frac{H_{\text{tmt}}(n, \Sigma[n]) = (\sigma', o, P') \qquad \Sigma' = \Sigma[n \mapsto \sigma']}{(P, \Sigma, T) \rightsquigarrow (P \uplus P', \Sigma', T + \langle \text{tmt}, o \rangle)} \text{ TIMEOUT}$$

Defining network semantics



Implementing Raft



Implementing Raft: Leader Election





Implementing Raft



Implementing Raft: Log Replication







Verification Challenge state machine replication

Raft Algorithm implemented in Verdi

Proof Overview and lessons learned







Verifying Raft: Show linearizability



Verifying Raft: Approach





State Machine Safety

Nodes agree about committed entries



since only committed entries executed



proof by induction on an execution

State Machine Safety: Proof



State Machine Safety: Proof



The burden of proof



Re-verification is the primary challenge:

- invariants are not inductive
- not-yet-verified code is wrong
- need additional invariants



The burden of proof



Re-verification is the primary challenge

Proof engineering techniques help:

- affinity lemmas
- intermediate reachability
- structural tactics
- information hiding

Ghost State: Example



Capture all entries received by a node



Affinity Lemmas: Example

 $e \in \log \implies$ e.term > 0 every invariant of entries in logs is invariant of entries in allEntries

 $e \in allEntries \implies e.term > 0$

Affinity Lemmas: Example

 $e \in \log \implies$ P e every invariant of entries in logs is invariant of entries in allEntries

$e \in allEntries \implies Pe$

Affinity Lemmas

Ex 1: Relate ghost state to real state transfer properties once and for all

Ex 2: Relate current messages to past response => past request







The burden of proof



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