

CSE 503

Software Engineering

Winter 2021

Intro to Abstract Interpretation

January 13, 2021

Recap: static vs. dynamic analysis

Static analysis

- Reason about the program without executing it.
- Build an abstraction of run-time states.
- Reason over abstract domain.
- Prove a property of the program.
- Sound* but conservative.

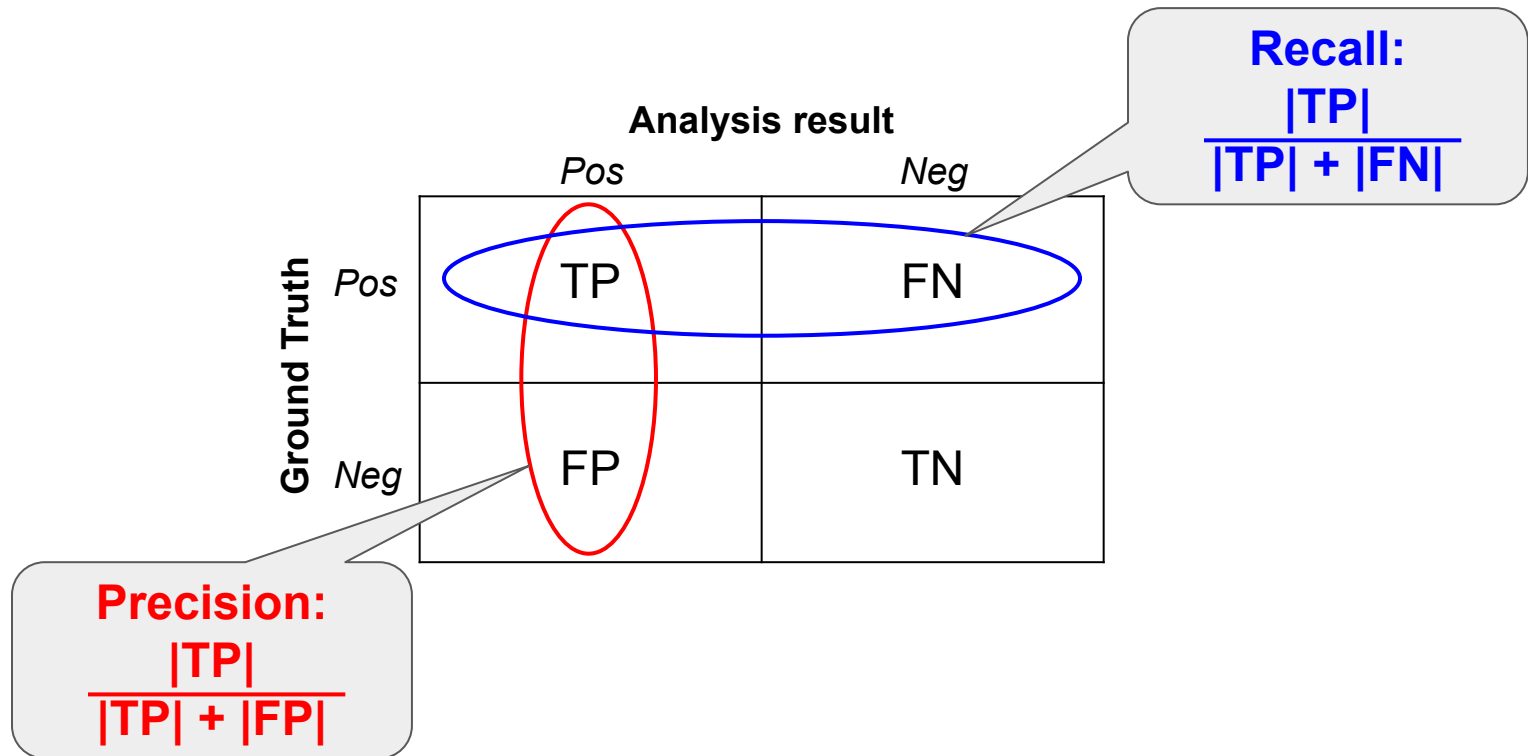
Dynamic analysis

- Reason about the program based on some program executions.
- Observe concrete behavior at run time.
- Improve confidence in correctness.
- Unsound* but precise.

* Some static analyses are unsound; dynamic analyses can be sound.

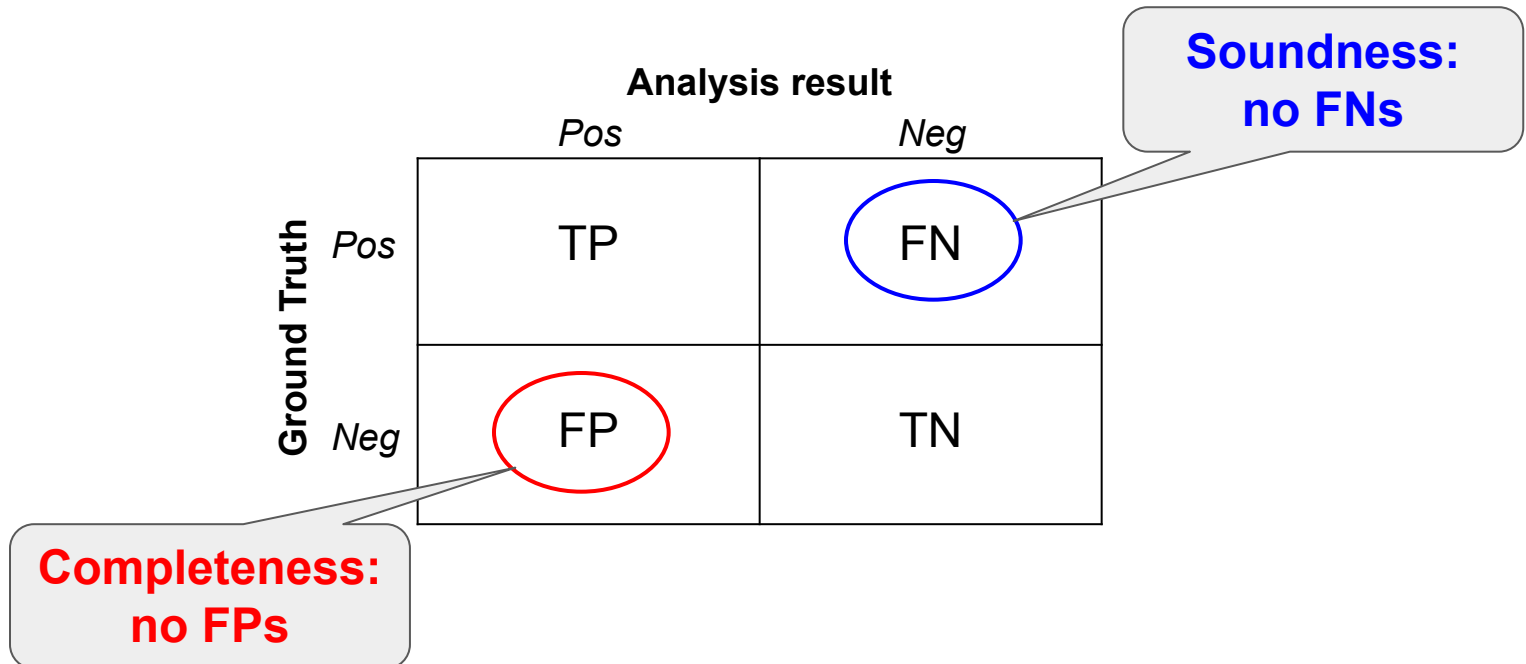
Recap: Terminology and important concepts

1. Precision vs. Recall (and FP/FN/TP/TN)



Recap: Terminology and important concepts

1. Precision vs. Recall (and FP/FN/TP/TN)
2. Soundness vs. Completeness



Recap: Terminology and important concepts

1. Precision vs. Recall (and FP/FN/TP/TN)
2. Soundness vs. Completeness
3. Concrete domain vs. Abstract domain

Concrete domain

0, 1, 2, 3, 4, ...

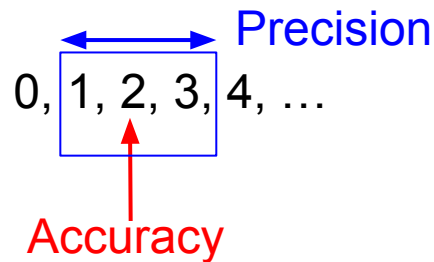
Abstract domain

even, odd

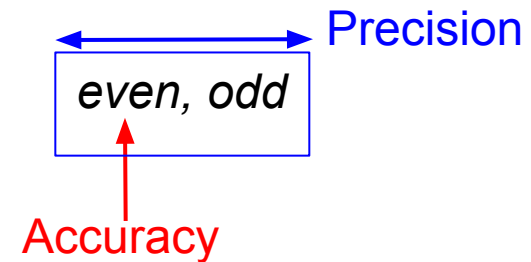
Recap: Terminology and important concepts

1. Precision vs. Recall (and FP/FN/TP/TN)
2. Soundness vs. Completeness
3. Concrete domain vs. Abstract domain
4. Accuracy vs. Precision

Concrete domain



Abstract domain

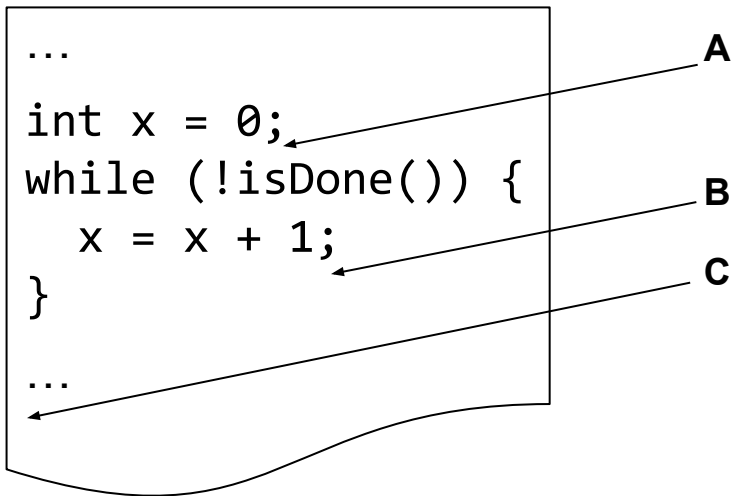


Today

- Abstract interpretation
 - Introduction
 - Abstraction functions
 - Concretization functions
 - Transfer functions
 - Lattices

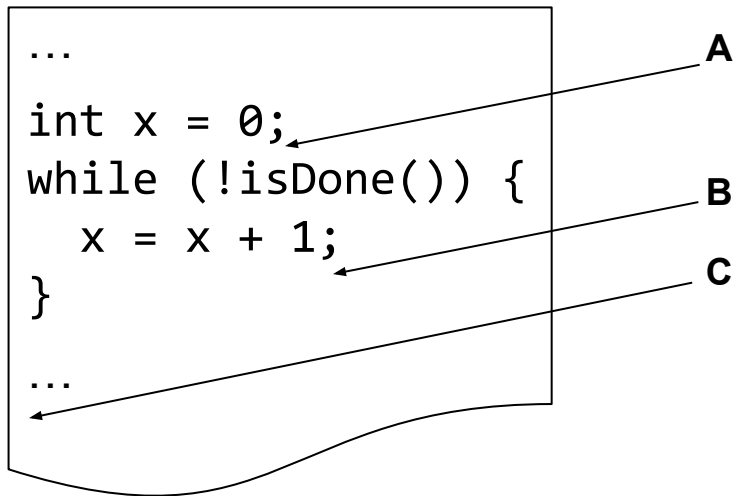
Properties of an ideal program analysis

- Soundness
- Completeness
- Termination



Properties of an ideal program analysis

- Soundness
- Completeness
- Termination



Abstract interpretation sacrifices completeness (precision)

Abstract interpretation: applications

Compiler checks and optimizations

- Liveness analysis (register reallocation)
- Reachability analysis (dead code elimination)
- Code motion (`while(cond){x = comp(); ...}`)

Abstract interpretation: code examples

Liveness

```
public class Liveness {  
    public void liveness() {  
        int a;  
        if (alwaysTrue()) {  
            a = 1;  
        }  
        System.out.println(a);  
    }  
}
```

Reachability

```
public void deadCode() {  
    return;  
    System.out.println("Here!");  
}
```

Abstract interpretation: example

Program

```
x = 0;  
y = read_even();  
x = y + 1;  
y = 2 * x;  
x = y - 2;  
y = x / 2;
```

Are all statements necessary?

Abstract interpretation: example

Program

```
x = 0;  
y = read_even();  
x = y + 1;  
y = 2 * x;  
x = y - 2;  
y = x / 2;
```

SSA form

```
x1 = 0;  
y1 = read_even();  
x2 = y1 + 1;  
y2 = 2 * x2;  
x3 = y2 - 2;  
y3 = x3 / 2;
```

x_1 is never read.

Abstract interpretation: example

Program

```
x = 0;  
y = read_even();  
x = y + 1;  
y = 2 * x;  
x = y - 2;  
y = x / 2;
```

SSA form

```
x1 = 0;  
y1 = read_even();  
x2 = y1 + 1;  
y2 = 2 * x2;  
x3 = y2 - 2;  
y3 = x3 / 2;
```

```
y3 = x3 / 2  
y3 = (y2 - 2) / 2  
y3 = (2 * x2 - 2) / 2  
y3 = (2 * (y1 + 1) - 2) / 2  
y3 = (2 * y1 + 2 - 2) / 2  
y3 = y1  
x3 = y1 * 2
```

Symbolic reasoning shows simplification potential.

Abstraction function

Program

```
x = 0;  
y = read_even();  
x = y + 1;  
y = 2 * x;  
x = y - 2;  
y = x / 2;
```

Concrete execution

```
{x=0; y=undef}  
{x=0; y=8}  
{x=9; y=8}  
{x=9; y=18}  
{x=16; y=18}  
{x=16; y=8}
```

Abstraction function

Program

```
x = 0;  
y = read_even();  
x = y + 1;  
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```

Mapping to abstract domain (even, odd)

```
—————> {x=e; y=e}  
—————> {x=o; y=e}  
—————> {x=o; y=e}  
—————> {x=e; y=e}  
—————> {x=e; y=e}
```


Abstraction function

Program

```
x = 0;  
y = read_even();  
x = y + 1;  
y = 2 * x;  
x = y - 2;  
y = x / 2;
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```

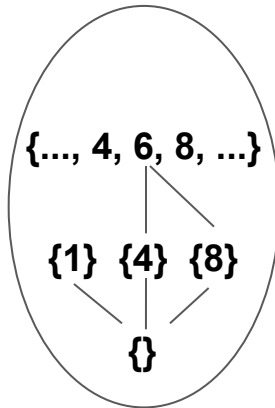
Mapping to abstract domain (even, odd)

```
—————→ {x=e; y=e}  
—————→ {x=o; y=e}  
—————→ {x=o; y=e}  
—————→ {x=e; y=e}  
—————→ {x=e; y=e}
```

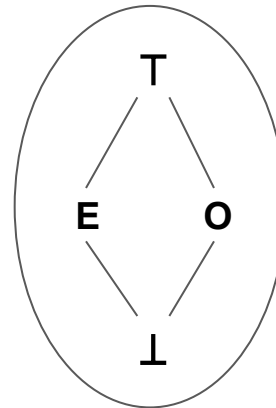
But, what's the purpose of the abstraction function?

Abstraction function

Concrete ($P(\mathbb{N})$)



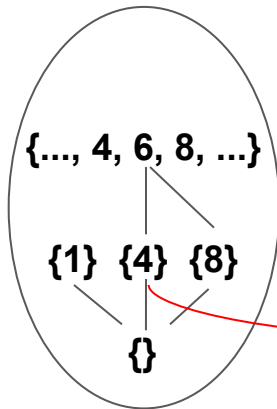
Abstract



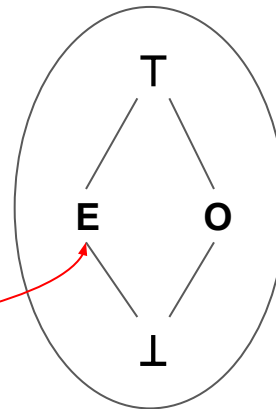
What is the abstraction (α) of $\{4\}$?

Abstraction function

Concrete ($P(\mathbb{N})$)



Abstract

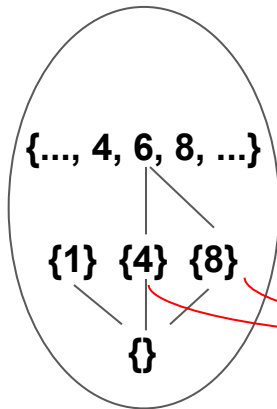


α

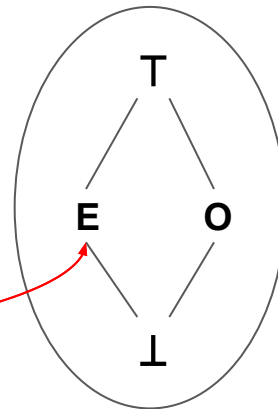
What is the abstraction (α) of $\{8\}$?

Abstraction function

Concrete ($P(\mathbb{N})$)



Abstract

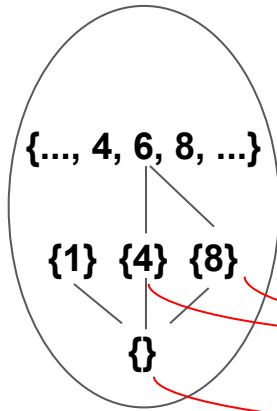


α

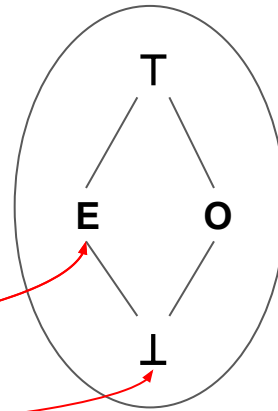
What is the abstraction (α) of $\{\}$?

Abstraction function

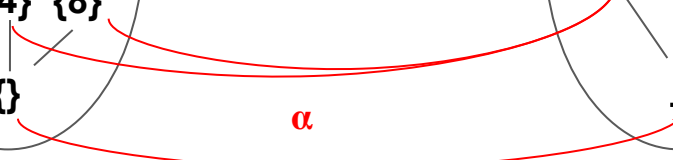
Concrete ($P(\mathbb{N})$)



Abstract



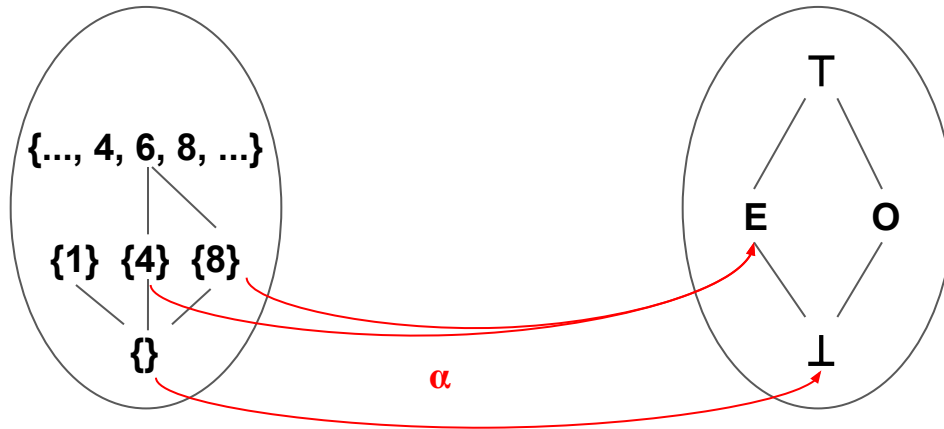
α



Concretization function

Concrete ($P(\mathbb{N})$)

Abstract

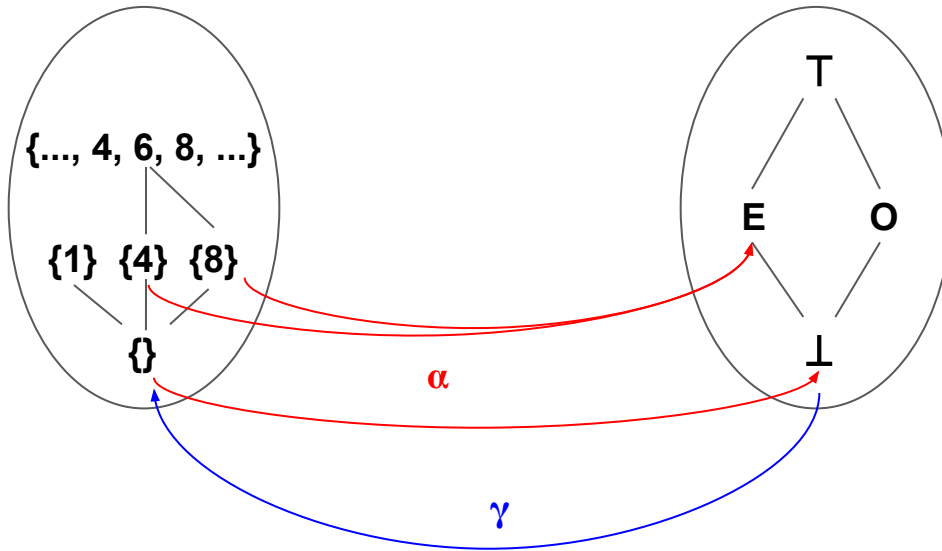


What is the concretization (γ) of \perp ?

Concretization function

Concrete ($P(\mathbb{N})$)

Abstract

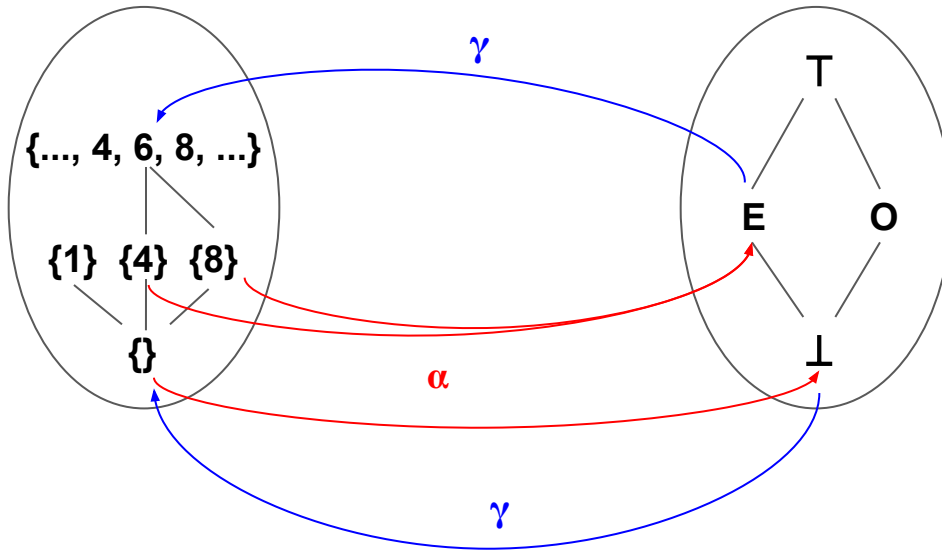


What is the concretization (γ) of E ?

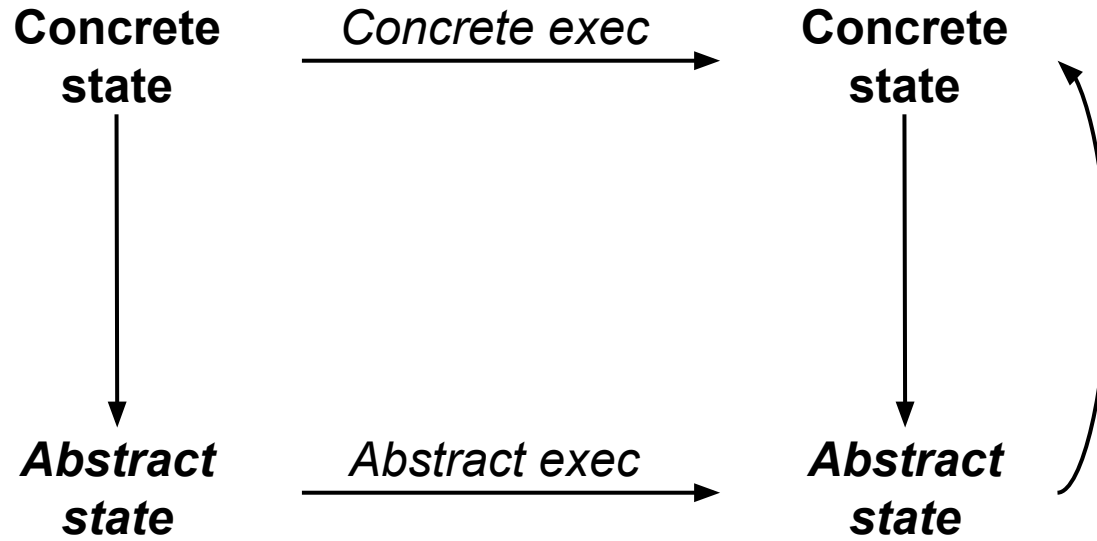
Concretization function

Concrete ($P(\mathbb{N})$)

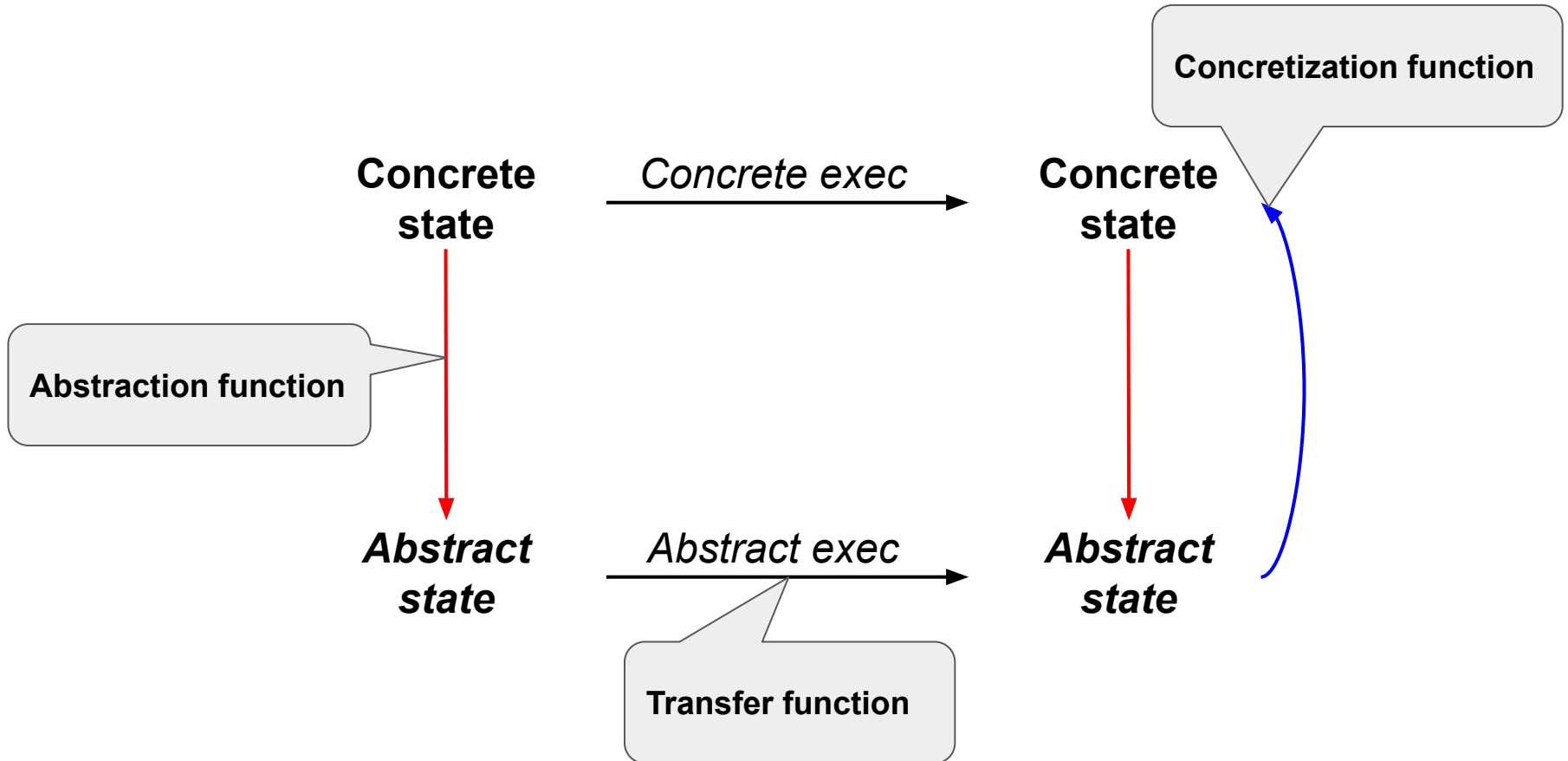
Abstract



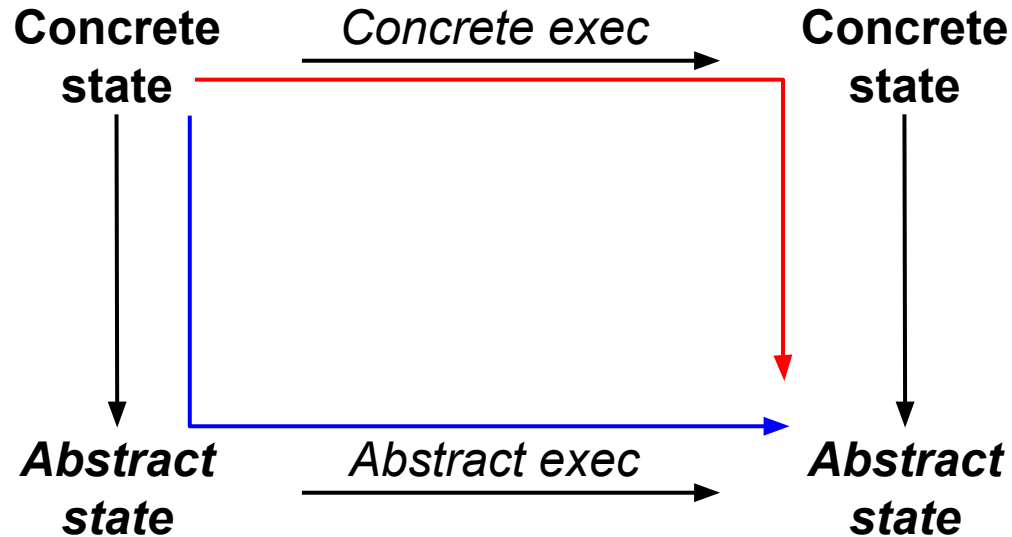
Transfer function



Transfer function

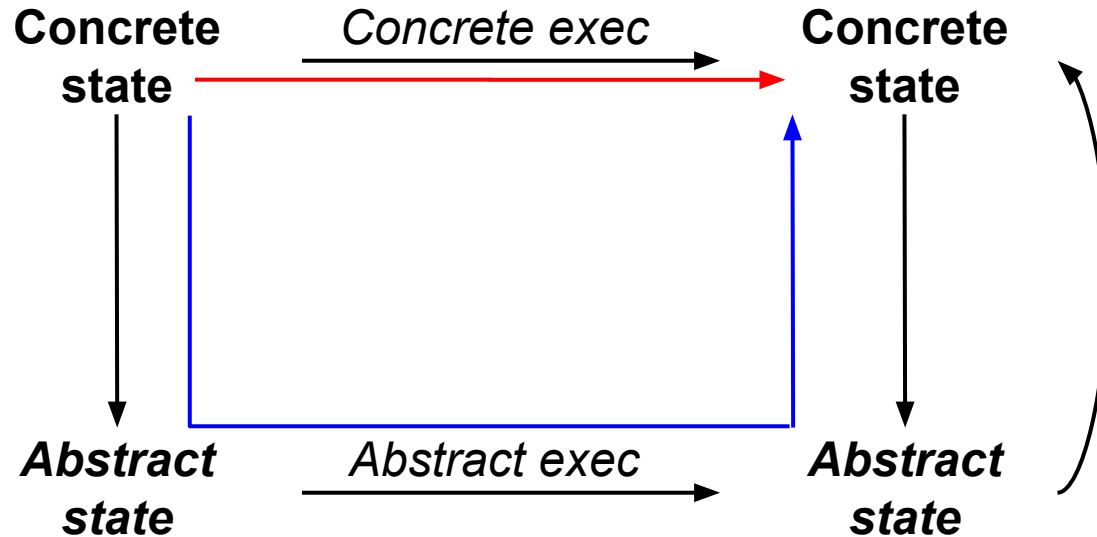


Abstract interpretation: approximation



Do both paths lead to the same abstract state?

Abstract interpretation: approximation



Do both paths lead to the same concrete state?

Set, semilattice, lattice

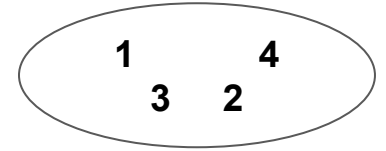
Set, semilattice, lattice

Set

Set, semilattice, lattice

Set

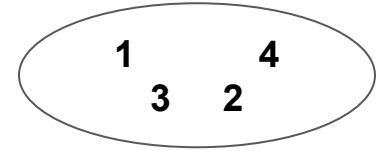
- **unordered** collection of **distinct** objects



Set, semilattice, lattice

Set

- unordered collection of distinct objects

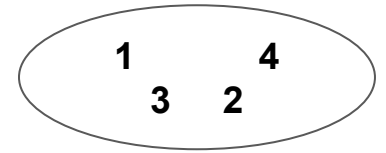


Partially ordered set

Set, semilattice, lattice

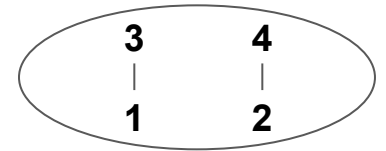
Set

- unordered collection of distinct objects



Partially ordered set

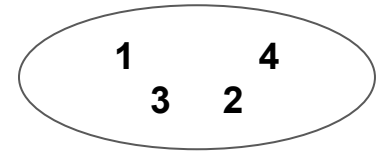
- Binary relationship \leq :
 - Reflexive: $x \leq x$
 - Anti-symmetric: $x \leq y \wedge y \leq x \Rightarrow x = y$
 - Transitive: $x \leq y \wedge y \leq z \Rightarrow x \leq z$



Set, semilattice, lattice

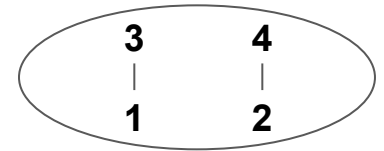
Set

- unordered collection of distinct objects



Partially ordered set

- Binary relationship \leq :
 - Reflexive: $x \leq x$
 - Anti-symmetric: $x \leq y \wedge y \leq x \Rightarrow x = y$
 - Transitive: $x \leq y \wedge y \leq z \Rightarrow x \leq z$



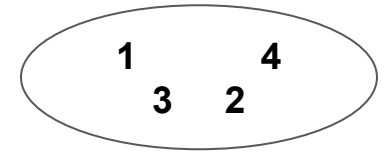
Join semilattice

Meet semilattice

Set, semilattice, lattice

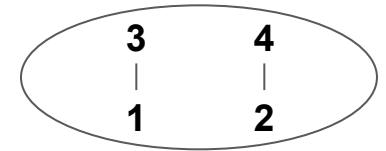
Set

- unordered collection of distinct objects



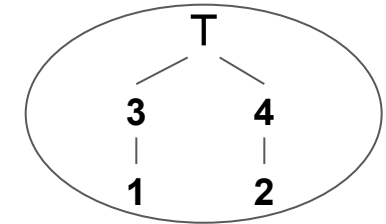
Partially ordered set

- Binary relationship \leq :
 - Reflexive: $x \leq x$
 - Anti-symmetric: $x \leq y \wedge y \leq x \Rightarrow x = y$
 - Transitive: $x \leq y \wedge y \leq z \Rightarrow x \leq z$



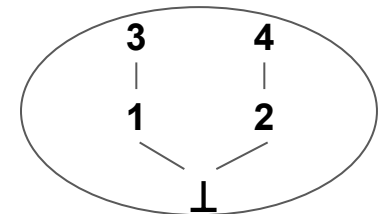
Join semilattice

- Partially ordered set with least upper bound (join)



Meet semilattice

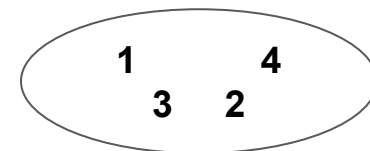
- Partially ordered set with greatest lower bound (meet)



Set, semilattice, lattice

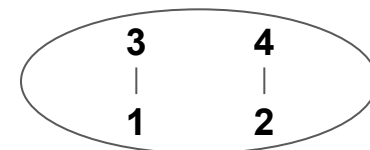
Set

- unordered collection of distinct objects



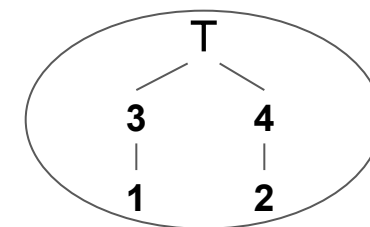
Partially ordered set

- Binary relationship \leq :
 - Reflexive: $x \leq x$
 - Anti-symmetric: $x \leq y \wedge y \leq x \Rightarrow x = y$
 - Transitive: $x \leq y \wedge y \leq z \Rightarrow x \leq z$



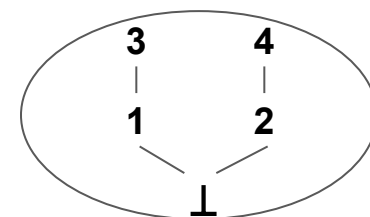
Join semilattice

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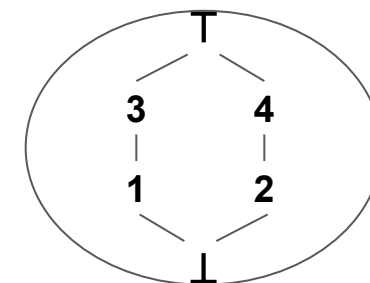
Meet semilattice

- Partially ordered set with greatest lower bound (meet)



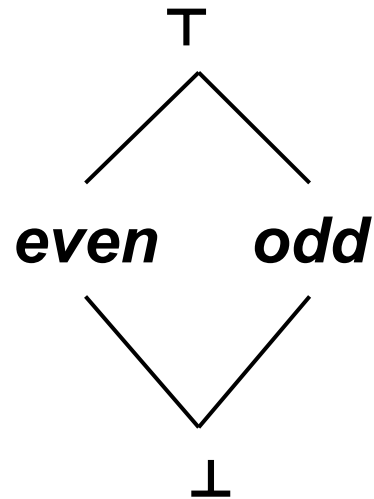
Lattice

- Both a join semilattice and a meet semilattice



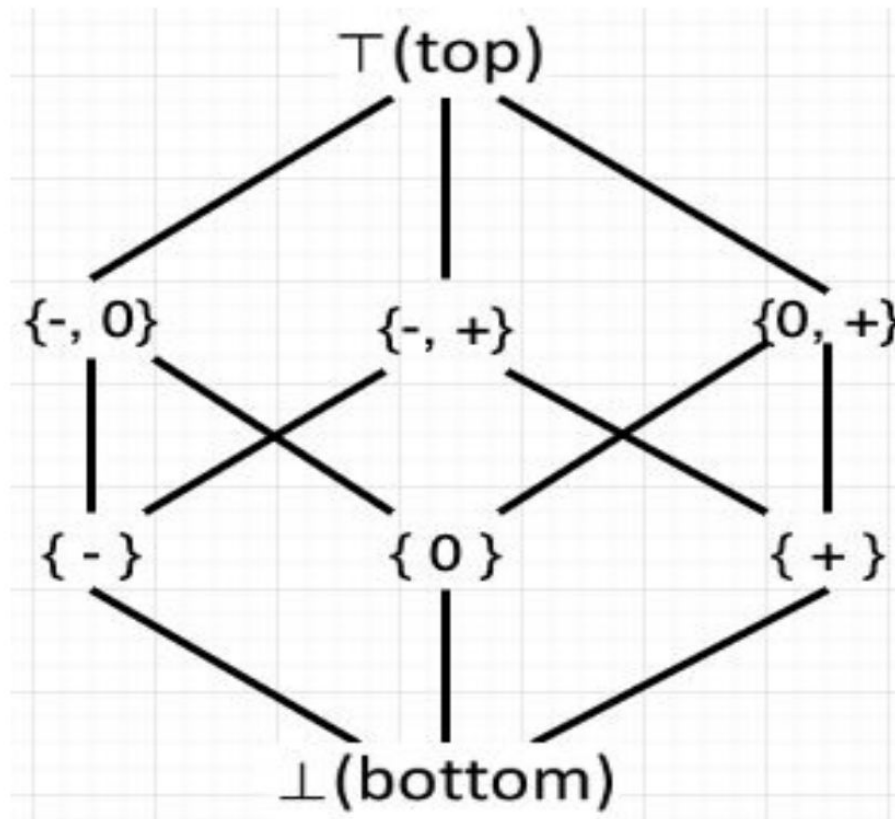
Lattice: example

Abstract domain: *even, odd, unknown* (\top), $\{\}$ (\perp)



Lattice: example

Abstract domain: $-$, 0 , $+$, *unknown*, $\{\}$



Lattice: example

Goal: approximate the values of x after the loop

```
...  
int x = 0;  
while (!isDone()) {  
    x = x + 1;  
}  
...
```

What are possible abstract domains and their trade-offs?

Lattice: example

Goal: approximate the values of x after the loop

```
...  
int x = 0;  
while (!isDone()) {  
    x = x + 1;  
}  
...
```

Possible abstract domains:

- Powerset of set of integers
- Intervals
- ...