Collaborative Verification of Information Flow for a High-Assurance App Store

Michael D. Ernst, **René Just**, Suzanne Millstein, Werner Dietl*, Stuart Pernsteiner, Franziska Roesner, Karl Koscher, Paulo Barros, Ravi Bhoraskar, Seungyeop Han, Paul Vines, and Edward X. Wu



University of Washington *University of Waterloo

November 6, 2014



Current commercial app stores



Current commercial app stores



Problem: Every major app store has approved malware!

Current commercial app stores



Problem: Every major app store has approved malware!

Best effort solution: Malware removed when encountered

High-assurance app stores

Needed in multiple domains

- Government app stores (e.g., DoD)
- Corporate app stores (e.g., financial sector)
- App stores for medical apps

Require stronger guarantees

Verified absence of (certain types of) malware

Verification is costly

- Effort is solely on app store side
- Analyst needs to understand/reverse-engineer the app

High-assurance app stores

Needed in multiple domains

- Government app stores (e.g., DoD)
- Corporate app stores (e.g., financial sector)
- App stores for medical apps

Require stronger guarantees

Verified absence of (certain types of) malware

Verification is costly

- Effort is solely on app store side
- Analyst needs to understand/reverse-engineer the app

Our solution: Collaboratively verify absence of malware Our focus: Information-flow malware

App

Permissions



Read location Internet

Арр

Permissions



Read location

Арр

Permissions



Read location



Read location Internet

Арр

Permissions



Read location



Camera



Information flow

Example: Information-flow malware



Read location

Permissions

Camera



Information flow

Example: Information-flow malware



Read location

Permissions

Camera





Permissions

Information flow

Read location Location \rightarrow Internet Internet

Location → Internet

Information flow













Approach: Overview

Collaborative verification model

Leverage but don't trust the developer

Information Flow Type-checker (IFT)

- Finer-grained permission model for Android
- False positives and declassifications
- Implicit information flow

Evaluation

- Effectiveness: Effective for real malware in real apps
- Usability: Low annotation and auditing burden









App store verifies



App store verifies

Developer and analyst do tasks that are easy for them

Verification of information flow



Verification of information flow



High-level description of permitted information flows

READ_SMS	->	INTERNET

- READ_CLIPBOARD -> DISPLAY
- USER_INPUT ->
- ACCESS_FINE_LOCATION ->
- - CALL PHONE
 - INTERNET(maps.google.com)

High-level description of permitted information flows

Source	flows to	Sink
	\sim	
READ_SMS	->	INTERNET
READ_CLIPBOARD	->	DISPLAY
USER_INPUT	->	CALL_PHONE
ACCESS_FINE_LOCATION	->	<pre>INTERNET(maps.google.com)</pre>

High-level description of permitted information flows

Source	flows to	Sink
READ_SMS	->	INTERNET
READ_CLIPBOARD	->	DISPLAY
USER_INPUT	->	CALL_PHONE
ACCESS_FINE_LOCATION	->	<pre>INTERNET(maps.google.com)</pre>

Sources and Sinks

Default Android permissions (145)

Not sufficient to model information flow!

High-level description of permitted information flows

Source	flows to	Sink
READ_SMS	->	INTERNET
READ_CLIPBOARD	->	DISPLAY
USER_INPUT	->	CALL_PHONE
ACCESS_FINE_LOCATION	->	<pre>INTERNET(maps.google.com)</pre>

Sources and Sinks

- Default Android permissions (145)
- Additional sensitive resources (28)

High-level description of permitted information flows

Source	flows to	Sink
	\sim	
READ_SMS	->	INTERNET
READ_CLIPBOARD	->	DISPLAY
USER_INPUT	->	CALL_PHONE
ACCESS_FINE_LOCATION	->	<pre>INTERNET(maps.google.com)</pre>

Sources and Sinks

- Default Android permissions (145)
- Additional sensitive resources (28)
- Parameterized permissions

Verification of information flow



Information flow types: Sources and Sinks

@Source Where might a value come from? **@Sink** Where might a value flow to?

Information flow types: Sources and Sinks

@Source Where might a value come from? **@Sink** Where might a value flow to?

Android API

void sendToInternet(String message);

String readGPS();

Information flow types: Sources and Sinks

@Source Where might a value come from? **@Sink** Where might a value flow to?


@Source Where might a value come from? **@Sink** Where might a value flow to?

Android API

void sendToInternet(@Sink(INTERNET)String message);

String readGPS();

@Source Where might a value come from? **@Sink** Where might a value flow to?



@Source Where might a value come from? **@Sink** Where might a value flow to?

Android API

void sendToInternet(@Sink(INTERNET)String message);

@Source(LOCATION)String readGPS();

@Source Where might a value come from? **@Sink** Where might a value flow to?

Android API

void sendToInternet(@Sink(INTERNET)String message);

```
@Source(LOCATION)String readGPS();
```

App code

String loc = readGPS();

sendToInternet(loc);

@Source Where might a value come from? **@Sink** Where might a value flow to?

Android API

void sendToInternet(@Sink(INTERNET)String message);

@Source(LOCATION)String readGPS();

App code

@Source(LOCATION)@Sink(INTERNET)String loc = readGPS(); sendToInternet(loc); Android API

Information flow types: Sources and Sinks

@Source Where might a value come from? **@Sink** Where might a value flow to?

API annotations are pre-verified

void sendToInternet(@Sink(INTERNET)String message);

@Source(LOCATION)String readGPS();

App codeDeveloper annotations are not trusted@Source(LOCATION)@Sink(INTERNET)String loc = readGPS();sendToInternet(loc);





@Source(ANY) = @Source({SMS, LOCATION, INTERNET, ...})











Verification of information flow



Information Flow Type-checker (IFT): Overview

Guarantees of type-checking

- 1. Annotations are consistent with code (type correctness)
- 2. Annotations are consistent with flow policy

Type checker verifies: annotations consistent

Information Flow Type-checker (IFT): Overview

Guarantees of type-checking

- 1. Annotations are consistent with code (type correctness)
- 2. Annotations are consistent with flow policy















```
App code

@Source({LOCATION, SMS})String[] array;

array[0] = readGPS();

array[1] = readSMS();

@Source(LOCATION)String loc = array[0];
```





```
App code
@Source({LOCATION, SMS})String[] array;
array[0] = readGPS();
array[1] = readSMS();
@SuppressWarnings("flow") // Always returns location data
@Source(LOCATION)String loc = array[0];
```

Declassifications

- Developer can suppress false-positive warnings
- App store employee verifies each declassification

Reducing false positives

Flow sensitivity

Type refinement with intra-procedural data flow analysis

```
App code
@Source({LOCATION, SMS})String value;
if (...) {
 value = readSMS();
 ... value: @Source(SMS)
}
... value: @Source({LOCATION, SMS})
```

Reducing false positives

Flow sensitivity

Type refinement with intra-procedural data flow analysis

Context sensitivity

Polymorphism (e.g., String operations, I/O streams, etc.)



Reducing false positives

Flow sensitivity

Type refinement with intra-procedural data flow analysis

Context sensitivity

Polymorphism (e.g., String operations, I/O streams, etc.)

Indirect control flow

- Constant value propagation
- Reflection analysis
- Intent analysis

App code

```
@Source(USER_INPUT)long creditCard = getCard();
long i=0;
while (true) {
    if (++i == creditCard) {
        sendToInternet(i);
    }
}
```



Classic approach (Denning and Denning, CACM'77)

- Taint all computations in dynamic scope
- Over-tainting may lead to taint explosion



Our approach: Prune irrelevant conditions

- Add additional Sink CONDITIONAL
- Type-checking warning for conditions with sensitive Source



Our approach: Prune irrelevant conditions

- Add additional Sink CONDITIONAL
- Type-checking warning for conditions with sensitive Source
- Analyst is aware of context
- No need to analyze dynamic scope for irrelevant conditions (e.g., null checks, malicious conditions, or trigger)

Evaluation: Overview

Are our permission model and type system effective?

- Adversarial red team challenge
- Evaluation of effectiveness for real malware

Is our approach effective and efficient in a timeconstrained set up?

- Control team study
- Comparison of effectiveness and efficiency to control team

Is our verification model applicable for real-world apps?

- Usability study with annotators and auditors
- Evaluation of annotation and auditing burden

Evaluation: Overview

Are our permission model and type system effective?

- Adversarial red team challenge
- Evaluation of effectiveness for real malware

Is our approach effective and efficient in a timeconstrained set up?

- Control team study
- Comparison of effectiveness and efficiency to control team

Is our verification model applicable for real-world apps?

- Usability study with annotators and auditors
- Evaluation of annotation and auditing burden

Apps are not pre-annotated

Adversarial red team challenge

Setup

- 5 independent red teams
- 72 Android apps (47 malicious information-flow malware)
- 8,000 LOC and 12 permissions on average
Adversarial red team challenge

Setup

- 5 independent red teams
- 72 Android apps (47 malicious information-flow malware)
- 8,000 LOC and 12 permissions on average

Results for 47 malicious apps



- Android permissions
- Additional Sources and Sinks
- Parameterized permissions
- Undetected
- 96% overall detection rate 4% require modeling of information flow paths (LOCATION -> ENCRYPT -> INTERNET)
- 60% of apps require our finer-grained sources and sinks

Control team study

Setup

- Control team using dynamic and static analysis tools
- 18 Android apps (13 malicious)
- 7,000 LOC and 16 permissions on average

Control team study

Setup

- Control team using dynamic and static analysis tools
- 18 Android apps (13 malicious)
- 7,000 LOC and 16 permissions on average



Usability study

Setup

- 2 groups acting as annotators and auditors
- 11 Android apps (1 malicious)
- 900 LOC and 12 permissions on average

Usability study

Setup

- 2 groups acting as annotators and auditors
- 11 Android apps (1 malicious)
- 900 LOC and 12 permissions on average

Annotation burden

- 96% of type annotations are inferred
- Annotations required: 6 per 100 lines of code
- Annotation time: 16 minutes per 100 lines of code

Most time spent on reverse engineering

Usability study cont.

Declassifications

- 50% of apps had no declassifications
- On average 3 declassification per 1,000 lines of code

IFT's features effectively reduce false positives

Usability study cont.

Declassifications

- 50% of apps had no declassifications
- On average 3 declassification per 1,000 lines of code

IFT's features effectively reduce false positives

Auditing burden

- Overall review time: 3 minutes per 100 lines of code
- ► 35% of time: review the flow policy
- 65% of time: review declassifications & conditionals

Only 23% of conditionals needed to be reviewed

Related work: Information flow

Jif (Myers, POPL'99)

- A security-typed language (incompatible Java extension)
- Supports dynamic checks and focuses on expressiveness

FlowDroid (Arzt et al., PLDI'14), SuSi (Rasthofer et al., NDSS'14)

- FlowDroid propagates sources and sinks found by SuSi
- SuSi classifies Android API methods using machine learning

IFT makes static verification of Android apps practical

- Finer-grained sources and sinks at type level
- Compiler plug-in using standard Java type annotations

Related work: Collaborative verification model

Verifying browser extensions

- ▶ IBEX (Guha et al., S&P'11)
 - Verification of Fine (ML dialect) against complex policies
- Lerner et al., ESORICS'13
 - Verification of private browsing using annotated JavaScript

IFT verifies information flow in Android apps using a high-level flow policy

Automated policy verification

- Crowd-sourcing (Agarwal & Hall, MobiSys'13)
- ► Natural language processing (Pandita et al., USENIX'13)
- Clustering (Gorla et al., ICSE'14).

Could aid manual verification of flow policies

Evaluation

Conclusions

Collaborative verification model

- Low overall verification effort for developer and app store analyst
- IFT combined with other analyses

Information Flow Type-checker (IFT)

- Context and flow-sensitive type system
- Fine-grained model for sources and sinks
- High-level information flow policy

Evaluation

- Detected 96% information-flow malware
- Low annotation and auditing burden
- Low false-positive rate



App store verifies





https://www.cs.washington.edu/sparta

René Just, UW CSE